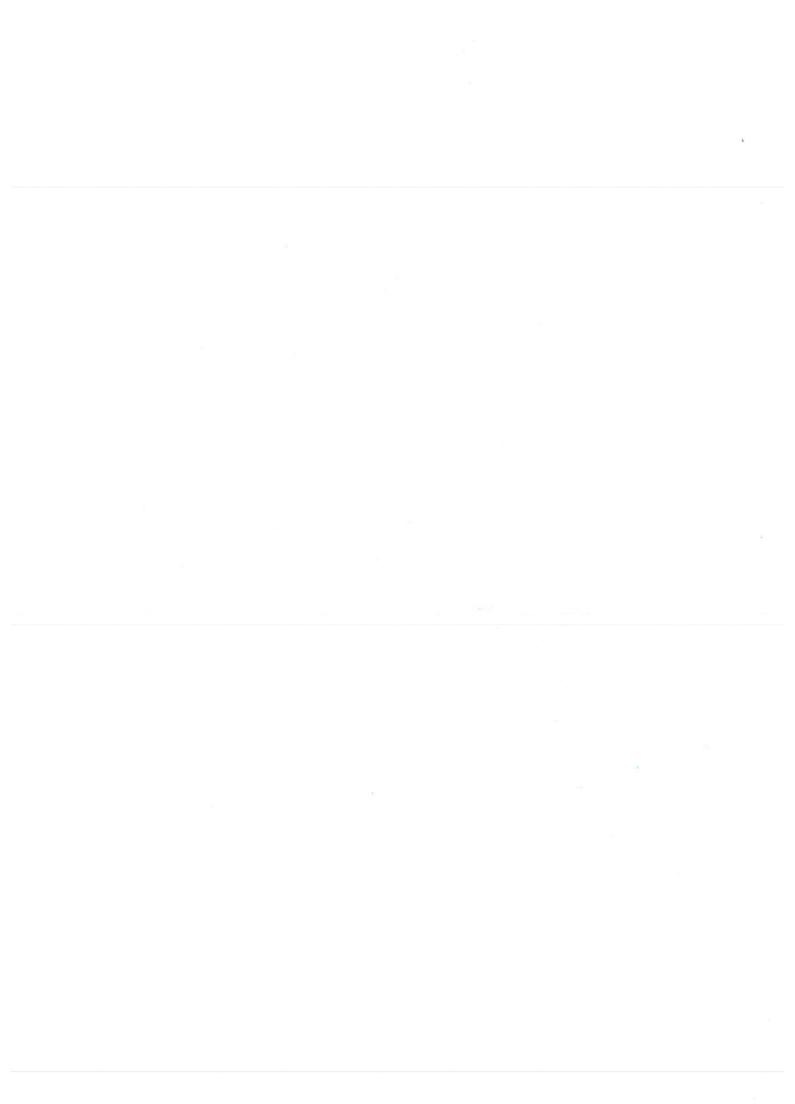
MATH0005 Algebra 1 Notes

Based on the 2018 autumn lectures by Prof F E A Johnson

The Author(s) has made every effort to copy down all the content on the board during lectures. The Author(s) accepts no responsibility for mistakes on the notes nor changes to the syllabus for the current year. The Author(s) highly recommends that the reader attends all lectures, making their own notes and to use this document as a reference only.

Recommended texts: () Heward Anton: Linear Algebra () Schawn Outline deries: Linear Algebra (S. Lipschutz) () Serge Lang Algebra.



Livier Eccarions : variables have exponent = 1.
$$\begin{cases} x_{ij} + z = 1 \\ x_{ij} + z = 1 \end{cases}$$

Tj we non out of symptons . $\begin{cases} w = x_{ij} + z = 1 \\ w = x_{ij} + z = 1 \end{cases}$
Tj we non out of symptons . $w = x_{ij} + z = 1 \end{cases}$
 $w = 1 \qquad w = 1$

MATRIX MULTIPLICATION

MXN Matrix

Define AB to be the mxp matrix. As gollows AB = (Cik) 1422m 16KEP. mxp matrix. Where Cik = (it row A). (Kth coulomn of B). example $B = \begin{pmatrix} 1 & 0 \\ 2 & -3 \\ 0 & 5 \end{pmatrix}$ $A = \begin{pmatrix} 1 & 2 & -1 \\ 0 & -2 & 3 \end{pmatrix}$ 2×3 3×2 AxB = 2xS matrix = $\begin{pmatrix} C_{11} & C_{12} \\ - & - & - \\ C_{21} & C_{22} \end{pmatrix}$ $C_{11} = \begin{pmatrix} \lambda^{bT} & row A \end{pmatrix} \cdot \begin{pmatrix} \lambda^{bT} & column & B \end{pmatrix} = \begin{pmatrix} \lambda, 2, -1 \end{pmatrix} \cdot \begin{pmatrix} \lambda \\ z \\ 4 \end{pmatrix} = \lambda + 4 - 4 = 4 \\ C_{12} = \begin{pmatrix} \lambda^{bT} & column & B \end{pmatrix} = \begin{pmatrix} \lambda, 2, -1 \end{pmatrix} \cdot \begin{pmatrix} 0 \\ -3 \\ -5 \end{pmatrix} = -6 - 5 = -11$ $C_{21} = (2^{nd} \operatorname{raw} A) \cdot (1^{st} \operatorname{column} B) = (0, -2, 3) \cdot (\frac{1}{2}) = 8$ $\begin{pmatrix} c_{22} = 0 + 6 + 15 = 21 \\ AB = \begin{pmatrix} 1 & -11 \\ 8 & 21 \end{pmatrix} \quad BA = \begin{pmatrix} 1 & 0 \\ 2 & -3 \\ 4 & 5 \end{pmatrix} \quad \begin{pmatrix} 1 & 2 & -1 \\ 0 & -2 & 3 \end{pmatrix} = \begin{pmatrix} 1 & 2 & -1 \\ 2 & 10 & -11 \\ 4 & -2 & 11 \end{pmatrix}$

 $AB \neq BA$

Formal definition $A = (a_{ij}) : \underline{Li} \underline{L} M \qquad B = (b_{jk}) : \underline{Lj} \underline{L} M \qquad 1 \leq \underline{Lj} \leq n$ $A = (a_{ij}) : \underline{Li} \underline{L} M \qquad B = (b_{jk}) : \underline{Lj} \leq n$ $A = (a_{ij}) : \underline{Li} \underline{L} M \qquad B = (b_{jk}) : \underline{Lj} \leq n$ $A = (a_{ij}) : \underline{Li} \underline{L} M \qquad B = (b_{ijk}) : \underline{Lj} \leq n$ $A = (b_{ijk}) : \underline{Lj} \leq n$ $A = (a_{ij}) : \underline{Li} \leq m$ $A = (b_{ijk}) : \underline{Lj} \leq n$ $A = (b_{ijk}) : \underline{Lj} \leq n$ $A = (b_{ijk}) : \underline{Lj} = (b_{ijk}) : \underline{$

$$\begin{aligned}
\mathcal{S} &= \begin{cases} a_{11} X_1 + \dots + a_{1N} X_n = b_1 \\ a_{21} X_1 + \dots + a_{2N} X_n = b_n \\ a_{21} X_1 + \dots + a_{2N} X_n = b_n \\ a_{21} X_1 + \dots + a_{2N} X_n = b_n \\ a_{21} X_1 + \dots + a_{2N} X_n = b_n \\ A_X = b_1 \quad \text{where} \\ A_X = b_1 \quad A_X = b_1 \\ A_X = b_1 \\ A_X = b_1 \\ A_X = b_1 \\ A_X =$$

It is true that : (AB) (= A (BC). Proof eventually.

· Distributive laws A is mxn A is $m \times n$ B, C $n \times p$. D $p \times q$ (B+C) = AB + AC (B+C) = BD + CD. · In general AB = BA Identity matrix In (each n 21). $\mathbf{I}_{2} = \begin{pmatrix} 4 & 0 \\ 0 & 4 \end{pmatrix}, \quad \mathbf{I}_{3} = \begin{pmatrix} 4 & 0 & 0 \\ 0 & 4 & 0 \\ 0 & 0 & 4 \end{pmatrix}, \quad \mathbf{I}_{4} = \begin{pmatrix} 4 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 4 & 0 \\ 0 & 0 & 0 & 4 \end{pmatrix}$ Konecker delta: $S_{ij} = \begin{cases} 1 & i=j \\ 0 & i\neq j \end{cases}$ $1 \leq i \leq n$ $S_{ij} = \begin{cases} 1 & i=j \\ 0 & i\neq j \end{cases}$ $1 \leq j \leq n$ $i \leq (2n)$ i = 2 0 = 1 0 = 1 0 = 1 0 = 1 0 = 1 0 = 1 0 = 1 0 = 1 0 = 1Definition: In= (Sij) 1620 - DIn = (Right Cold Cos nos con subindices In is identity for matine multiplication: (Right Cold Cos nos con subindices * Property A. (A. Multiplication: (Right Cold Cos nos con subindices is a low matine multiplication: (Right Cold Cos nos con subindices is a low matine multiplication: (Right Cold Cos nos con subindices is a low matine multiplication: (Right Cold Cos nos con subindices is a low matine multiplication: (Right Cold Cos nos con subindices is a low matine multiplication: (Right Cold Cos nos con subindices is a low matine multiplication: (Right Cold Cos nos con subindices is a low matine multiplication (Right Cold Cos nos con subindices is a low matine multiplication (Right Cold Cos nos con subindices is a low matine multiplication (Right Cold Cos nos con subindices is a low matine multiplication (Right Cold Cos nos con subindices is a low matine multiplication (Right Cold Cos nos con subindices is a low matine multiplication (Right Cold Cos nos con subindices). n is identity for mattice multipluce. K Reports $A = (a_{ij})_{ijk} \leq m$ $M \times a$ $a = (a_{ij})_{ijk} \leq a = (a_{ijk})_{ijk} < a = (a$ = $1.a_{ik+1} O \cdot S_{ij} = a_{ik} (SmA)_{ik} = a_{ik} : Ima = a_{ik}$ Deg? Let A = (aij) 1 Lich Say that A is invertible AB=In=BA A-1 A-1 A-1 A-1 14141 when there exists B (also nxn) s.t In ordinary aritmetric y a to there exists a = 1; ; aa'=1=a'a FALSE FOR MATRICES No sirve solo con que Azo, come venos en el ejemplo, aunque A = c puede no existic inverso

$$\frac{4}{2} \operatorname{respec}_{k} \left(\begin{array}{c} 0 \\ 0 \\ \end{array} \right) = \left(\begin{array}{c} 0 \\ 0 \\ \end{array} \right) \left(\begin{array}{c} 0 \end{array}$$

• Three types of operators
(a) Add
$$\lambda \in q_{1}(j)$$
 to $eq_{1}(\lambda) = E(\lambda,j,\lambda)$ - have of operators
(b) multiply experient(λ by $\lambda \neq 0 = iN(\lambda, \lambda)$ - have of operators
(c) multiply experient(λ by $\lambda \neq 0 = iN(\lambda, \lambda)$ - have of operators
(c) Indicatence over $q_{2} \in q_{1}(\lambda)$ and $\lambda \in u_{2}(j) = S(\lambda, j)$ - have of examination
(c) Indicatence over $q_{2} \in q_{1}(\lambda)$, $\lambda \neq 0$
(c) $M = (\lambda, \lambda)$; $M = (\lambda, \lambda)$; $A \neq 0$
(c) $M = (\lambda, \lambda)$; $M = (\lambda, \lambda)$; $A \neq 0$
(c) $M = (\lambda, \lambda)$;
(c) $A = \begin{pmatrix} a & b \\ c & d \\ d = d \\$

But
$$\int_{0}^{\infty} = 0$$
 when $5 \neq j$; $\int_{0}^{\infty} = 1$.
 $I \in (i, j) \land J_{i+1} = \int_{0}^{\infty} (\Lambda_{0}^{-1} - \frac{1}{2}) \land J_{0}^{-1} = J \land J_{0}^{-1} \land J_{0$

$$E(i_{1,j}, \lambda) = Sn+ \lambda E(i_{1,j}) \left(\begin{smallmatrix} 4460\\ 0 & 0 \end{smallmatrix} \right) \left(\begin{smallmatrix} 600\\ 2 & 00 \end{smallmatrix} \right) = \left[E(i_{1,j}, \lambda) = Sn+ \lambda E(i_{1,j}) \right] = \left[F(i_{1,j}, \lambda) = F(i_{1,j}, \lambda) = F(i_{1,j}, \lambda) \right] = \left[F(i_{1,j}, \lambda) = F(i_{1,j},$$

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Octobes 14th 2018

$$\begin{array}{l} \underbrace{\mathbb{E}(i,j,\lambda)}_{\substack{(i,j,\lambda) \in \mathbb{R}^{n} \\ \in \{i,j,\lambda\}}} & \underbrace{\mathbb{E}(i,j,\lambda)}_{\substack{(i,j,\lambda) \in \mathbb{R}^{n} \\ \in \{i,j,\lambda\}}} \underbrace{\mathbb{E}(i,j,\lambda)}_{\substack{(i,j,\lambda) \in \mathbb{R}^{n} \\ \in \{i,j,\lambda\}}} \underbrace{\mathbb{E}(i,j,\lambda)}_{\substack{(i,j,\lambda) \in \mathbb{R}^{n} \\ \in \mathbb{R}^{n} \\ \in \mathbb{R}^{n} \\ \in \mathbb{R}^{n} \\ \underbrace{\mathbb{E}(i,j,\lambda)}_{\substack{(i,j,\lambda) \in \mathbb{R}^{n} \\ \in \mathbb{R}^{n} \\ \in \mathbb{R}^{n} \\ \in \mathbb{R}^{n} \\ \underbrace{\mathbb{E}(i,j,\lambda)}_{\substack{(i,j,\lambda) \in \mathbb{R}^{n} \\ \in \mathbb{R}^{n} \\ \in \mathbb{R}^{n} \\ \in \mathbb{R}^{n} \\ \underbrace{\mathbb{E}(i,j,\lambda)}_{\substack{(i,j,\lambda) \in \mathbb{R}^{n} \\ \in \mathbb{R}^{n} \\ \in \mathbb{R}^{n} \\ \in \mathbb{R}^{n} \\ \underbrace{\mathbb{E}(i,j,\lambda)}_{\substack{(i,j,\lambda) \in \mathbb{R}^{n} \\ \in \mathbb{R}^{n} \\ \in \mathbb{R}^{n} \\ \in \mathbb{R}^{n} \\ \underbrace{\mathbb{E}(i,j,\lambda)}_{\substack{(i,j,\lambda) \in \mathbb{R}^{n} \\ \underbrace{\mathbb{E}(i,j,\lambda)}_{\substack{(i,j,\lambda) \in \mathbb{R}^{n} \\ \in \mathbb{R}^{n} \\ \in \mathbb{R}^{n} \\ \in \mathbb{R}^{n} \\ \underbrace{\mathbb{E}(i,j,\lambda)}_{\substack{(i,j,\lambda) \in \mathbb{R}^{n} \\ \underbrace{\mathbb{E}(i,j,\lambda)}_{\substack{(i,j,\lambda) \in \mathbb{R}^{n} \\ \underbrace{\mathbb{E}(i,j,\lambda)}_{\substack{(i,j,\lambda) \in \mathbb{R}^{n} \\ \underbrace{\mathbb{E}(i,\lambda)}_{\substack{(i,j,\lambda) \in \mathbb{R}^{n} \\ \underbrace{\mathbb{E}(i,\lambda)}_{\substack{(i,j,\lambda) \in \mathbb{R}^{n} \\ \in \mathbb$$

$$\frac{\operatorname{Comple}_{(i)} \operatorname{m=4}_{(i)} \cdots \operatorname{m=4}_{(i)} \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{array} \right) \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 \end{array} \right) \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{array} \right) \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{array} \right) \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{array} \right) \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{array} \right) \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 \\ 0 & 0 & 0 & 0 \end{array} \right) \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 \\ 0 & 0 & 0 & 0 \end{array} \right) \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 \\ 0 & 0 & 0 & 0 \end{array} \right) \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 \\ 0 & 0 & 0 & 0 \end{array} \right) \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 \\ 0 & 0 & 0 & 0 \end{array} \right) \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 \\ 0 & 0 & 0 & 0 \end{array} \right) \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 \\ 0 & 0 & 0 & 0 \end{array} \right) \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 \\ 0 & 0 & 0 & 0 \end{array} \right) \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 \\ 0 & 0 & 0 & 0 \end{array} \right) \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 \\ 0 & 0 & 0 & 0 \end{array} \right) \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 \\ 0 & 0 & 0 & 0 \end{array} \right) \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 \\ 0 & 0 & 0 & 0 \end{array} \right) \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 \\ 0 & 0 & 0 & 0 \end{array} \right) \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 \\ 0 & 0 & 0 & 0 \end{array} \right) \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 \\ 0 & 0 & 0 & 0 \end{array} \right) \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 \\ 0 & 0 & 0 & 0 \end{array} \right) \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 \\ 0 & 0 & 0 & 0 \end{array} \right) \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 \\ 0 & 0 & 0 & 0 \end{array} \right) \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 \\ 0 & 0 & 0 & 0 \end{array} \right) \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 \\ 0 & 0 & 0 & 0 \end{array} \right) \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 \\ 0 & 0 & 0 & 0 \end{array} \right) \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 \\ 0 & 0 & 0 & 0 \end{array} \right) \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 \\ 0 & 0 & 0 \end{array} \right) \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 \\ 0 & 0 & 0 \end{array} \right) \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 \\ 0 & 0 & 0 \end{array} \right) \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 \\ 0 & 0 & 0 \end{array} \right) \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 \\ 0 & 0 & 0 \end{array} \right) \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 \\ 0 & 0 & 0 \end{array} \right) \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 \\ 0 & 0 & 0 \end{array} \right) \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 \\ 0 & 0 & 0 \end{array} \right) \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 \\ 0 & 0 & 0 \end{array} \right) \cdots \operatorname{m=1}_{(i)} \left(\begin{array}{c} a & 0 \\ \cdots \end{array})$$

Formal def of
$$P(i,j)$$
.
 $P(i,j) = Im - E(i,i) - E(j,j) + E(i,j) + E(j,i)$.

SYSTEMS OF OPERATIONS .

Certain systems of equations have obvious solutions.

$$\frac{\text{Example}}{\left(\begin{array}{c} x_{1} \\ x_{2} \end{array}\right)} + \frac{x_{2}}{\left(\begin{array}{c} x_{3} \\ x_{3} \end{array}\right)} + \frac{x_{4}}{\left(\begin{array}{c} x_{5} \\ x_{5} \end{array}\right)} - \frac{x_{6}}{\left(\begin{array}{c} x_{5} \\ x_{5} } - \frac{x_{6}}{\left(\begin{array}{c} x_{5} } - \frac{x_{6}}{\left(\begin{array}{c} x_{5} \\ x_{5} } - \frac{x_{6}}{\left(\begin{array}{c} x_{5} \\ x_{5} } - \frac{x_{6}}{\left(\begin{array}{c} x_{5} \\ x_{5}$$

 $X_{1} = \lambda = X_{2} + X_{4} - 2X_{6}$ $X_{2} = X_{2}$ $X_{3} = 2 - X_{4} - X_{6}$ $X_{7} = X_{4}$ $X_{5} = 3 + X_{6}$ $X_{6} = X_{6}$ $X_{6} = X_{6}$ $X_{1} = \lambda_{1}$ $X_{2} = \lambda_{1}$ $X_{3} = \lambda_{1}$ $X_{4} = \lambda_{1}$ $X_{5} = \lambda_{1}$ $X_{6} = \lambda_{1}$

 $\begin{array}{c} (X_{1}) \\ (X_{2}) \\ (X_{3}) \\ (X_{3})$

(Reduced) Row Echelon Matix: 1) In any nontern row 1st element you meet must be 1 (leading 1) 2) The rest of the column of a leading 1 must be 0. 3) Rows are stepped in Echelon form. 4) Any two rows notes and the row may record row.

$$\begin{array}{l} \underbrace{\operatorname{Ecompter}_{i}}_{\substack{(1 \leq i \leq i \leq i \leq i \\ 0 \leq i \leq i \leq i \leq i \\ 0 \leq i \leq i \leq i \leq i \\ 0 \leq i \leq i \leq i \leq i \\ 0 \leq i \leq i \leq i \leq i \\ 0 \leq i \leq i \leq i \leq i \\ \hline \begin{pmatrix} 1 \leq i \leq i \leq i \leq i \\ 0 \leq i \leq i \leq i \\ 0 \leq i \leq i \leq i \\ 0 \leq i \leq i \leq i \leq i \\ \hline \begin{pmatrix} 1 \leq i \leq i \leq i \leq i \\ 0 \leq i \leq i \leq i \\ 0 \leq i \leq i \leq i \\ 0 \leq i \leq i \leq i \\ \hline \end{pmatrix} \rightarrow \operatorname{It} i_{0} \quad \operatorname{rot} \text{ stepped} \left\{ \begin{array}{c} 11 \\ 1 \\ 1 \\ 0 \leq i \leq i \leq i \\ \hline \\ 0 \leq i \leq i \leq i \\ \hline \\ 0 \leq i \leq i \leq i \\ \hline \\ 0 \leq i \leq i \leq i \\ \hline \\ 0 \leq i \leq i \leq i \\ \hline \\ 0 \leq i \\ \hline \\$$

$$\begin{split} & \underbrace{\mathsf{Condit}}_{\text{Conditional}} & \underbrace{\mathsf{N}_{1} - \mathsf{X}_{2} + \mathsf{X}_{3} - \mathsf{X}_{4} + \mathsf{X}_{7} - \mathsf{X}_{6} = 2}_{X_{1} - \mathsf{X}_{2} + 2\mathsf{X}_{3} - \mathsf{X}_{7} - \mathsf{X}_{6} = 2}_{X_{1} - \mathsf{X}_{2} - 2}_{X_{1} - \mathsf{X}_{2} - 2}_{X_{1} - \mathsf{X}_{3} - \mathsf{X}_{5} - \mathsf{X}_{6} = 0}_{X_{1} - \mathsf{X}_{2} - 2}_{X_{1} - \mathsf{X}_{2} - 2}_{X_{1} - \mathsf{X}_{3} - \mathsf{X}_{5} - \mathsf{X}_{6} = 0}_{X_{1} - \mathsf{X}_{2} - \mathsf{X}_{2} - \mathsf{X}_{3} - \mathsf{X}_{5} = 0}_{X_{1} - \mathsf{X}_{2} - \mathsf{X}_{3} - \mathsf{X}_{5} - \mathsf{X}_{6} = 0}_{X_{1} - \mathsf{X}_{2} - \mathsf{X}_{3} - \mathsf{X}_{5} - \mathsf{X}_{6} = 0}_{X_{1} - \mathsf{X}_{2} - \mathsf{X}_{3} - \mathsf{X}_{5} - \mathsf{X}_{6} = 0}_{X_{1} - \mathsf{X}_{2} - \mathsf{X}_{3} - \mathsf{X}_{3} - \mathsf{X}_{5}}_{X_{2}} \\ & \underbrace{\mathsf{Preme augmentates matrix}_{1 - \mathsf{X}_{2} - \mathsf{X}_{3} - \mathsf{X}_{5}}_{X_{2} - \mathsf{X}_{1} - \mathsf{X}_{2} - \mathsf{X}_{2}}_{X_{1} - \mathsf{X}_{2} - \mathsf{X}_{2}}_{X_{2} - \mathsf{X}_{2} - \mathsf{X}_{5} - \mathsf{X}_{6}}_{X_{2} - \mathsf{X}_{3}}_{X_{3} - \mathsf{X}_{4}}_{X_{3} - \mathsf{X}_$$

(2)
$$F_{z} = F_{z} = F_{z} + F_{z} + F_{z} - F_{y} + F_{z} = F_{z} + F_{z} +$$

write down the augmented matrix:

$$(A1b) = \begin{pmatrix} 1 & 1 & -1 & -1 & | & 4 \\ 1 & -1 & 1 & -1 & | & -1 & | & 5 \\ 1 & 5 & 1 & -1 & 1 & -1 & | & 5 \\ 1 & 5 & 1 & -1 & 1 & -1 & | & 2 \end{pmatrix}$$

Perform operations to augmented matrix in to reduced Row Echelon For

$$\begin{pmatrix} | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & | 1 & |$$

$$\begin{vmatrix} 9/2 & -X_3 + X_7 - X_5 + X_6 \\ -1/2 \\ X_3 \\ X_4 \\ X_5 \\ X_6 \end{vmatrix}$$
 so many collutions $X_3, X_4, 55, K_6$
arbitrary.

$$\begin{split} & \left\{ \begin{array}{c} \left\{ \begin{array}{c} x_{1} + x_{2} &= 1 \\ x_{1} + x_{2} + x_{3} &= 5 \\ x_{2} + x_{3} &= 4 \end{array} \right. \\ & \left\{ \begin{array}{c} \left\{ \begin{array}{c} 1 & 1 & 0 \\ 1 & 1 & 1 \\ 0 & 1 & 1 \end{array} \right\} \\ \left\{ \begin{array}{c} \left\{ \begin{array}{c} 1 & 1 & 0 \\ 0 & 1 & 1 \end{array} \right\} \\ \left\{ \begin{array}{c} 1 & 1 \\ 0 & 1 & 1 \end{array} \right\} \\ \left\{ \begin{array}{c} 1 & 1 \\ 0 & 1 \end{array} \right\} \\ \left\{ \begin{array}{c} 1 & 1 \\ 0 & 1 \end{array} \right\} \\ \left\{ \begin{array}{c} 1 & 1 \\ 0 & 1 \end{array} \right\} \\ \left\{ \begin{array}{c} 1 & 1 \\ 0 & 1 \end{array} \right\} \\ \left\{ \begin{array}{c} 1 & 1 \\ 0 & 1 \end{array} \right\} \\ \left\{ \begin{array}{c} 1 & 1 \\ 0 & 1 \end{array} \right\} \\ \left\{ \begin{array}{c} 1 & 1 \\ 0 & 1 \end{array} \right\} \\ \left\{ \begin{array}{c} 1 & 1 \\ 0 & 1 \end{array} \right\} \\ \left\{ \begin{array}{c} 1 & 1 \\ 0 & 1 \end{array} \right\} \\ \left\{ \begin{array}{c} 1 & 1 \\ 0 & 1 \end{array} \right\} \\ \left\{ \begin{array}{c} 1 & 1 \\ 0 & 1 \end{array} \right\} \\ \left\{ \begin{array}{c} 1 & 1 \\ 0 & 1 \end{array} \right\} \\ \left\{ \begin{array}{c} 1 & 1 \\ 0 \end{array} \right\} \\ \left\{ \begin{array}{c} 1 & 0 \\ 0 & 0 \end{array} \right\} \\ \left\{ \begin{array}{c} 1 & 1 \\ 0 \end{array} \right\} \\ \left\{ \begin{array}{c} 1 & 0 \\ 0 \end{array} \right\} \\ \left\{ \begin{array}{c} 1 & 0 \\ 0 \end{array} \right\} \\ \left\{ \begin{array}{c} 1 & 0 \\ 0 \end{array} \right\} \\ \left\{ \begin{array}{c} 1 & 0 \\ 0 \end{array} \right\} \\ \left\{ \begin{array}{c} 1 & 0 \\ 0 \end{array} \right\} \\ \left\{ \begin{array}{c} 1 & 0 \\ 0 \end{array} \right\} \\ \left\{ \begin{array}{c} 1 & 0 \\ 0 \end{array} \right\} \\ \left\{ \begin{array}{c} 1 & 0 \\ 0 \end{array} \right\} \\ \left\{ \begin{array}{c} 1 & 0 \\ 0 \end{array} \right\} \\ \left\{ \begin{array}{c} 1 & 0 \\ 0 \end{array} \right\} \\ \left\{ \begin{array}{c} 1 & 0 \\ 0 \end{array} \right\} \\ \left\{ \begin{array}{c} 1 & 0 \\ 0 \end{array} \right\} \\ \left\{ \begin{array}{c} 1 & 0 \\ 0 \end{array} \right\} \\ \left\{ \begin{array}{c} 1 & 0 \\ 0 \end{array} \right\} \\ \left\{ \begin{array}{c} 1 & 0 \\ 0 \end{array} \right\} \\ \left\{ \begin{array}{c} 1 & 0 \\ 0 \end{array} \\ \left\{ \begin{array}{c} 1 & 0 \\ 0 \end{array} \right\} \\ \left\{ \begin{array}{c} 1 & 0 \\ 0 \end{array} \\ \left\{ \begin{array}{c} 1 & 0 \\ 0 \end{array} \\ \left\{ \begin{array}{c} 1 & 0 \\ 0 \end{array} \\ \left\{ \begin{array}{c} 1 & 0 \\ 0 \end{array} \\ \left\{ \begin{array}{c} 1 & 0 \end{array} \\ \left\{ \begin{array}{c} 1 & 0 \\ 0 \end{array} \\ \left\{ \begin{array}{c} 1 & 0 \end{array} \\ \left\{ \begin{array}{c} 1 & 0 \end{array} \\ 0 \end{array} \\ \left\{ \begin{array}{c} 1 & 0 \end{array} \\ \left\{ \begin{array}{c} 1 & 0 \end{array} \\ \left\{ \begin{array}{c} 1 & 0 \end{array} \\ 0 \end{array} \\ \left\{ \begin{array}{c} 1 & 0 \end{array} \\ \left\{ \begin{array}{c} 1 & 0 \end{array} \\ 0 \end{array} \\ \left\{ \begin{array}{c} 1 & 0 \end{array} \\ \left\{ \begin{array}{c} 1 & 0 \end{array} \\ 0 \end{array} \\ \left\{ \begin{array}{c} 1 & 0 \end{array} \\ 0 \end{array} \\ \left\{ \begin{array}{c} 1 & 0 \end{array} \\ 0 \end{array} \\ \left\{ \begin{array}{c} 1 & 0 \end{array} \\ 0 \end{array} \\ \left\{ \begin{array}{c} 1 & 0 \end{array} \\ \left\{ \begin{array}{c} 1 & 0 \end{array} \\ 0 \end{array} \\ \left\{ \begin{array}{c} 1 & 0 \end{array} \\ 0 \end{array} \\ \left\{ \begin{array}{c} 1 & 0 \end{array} \\ 0 \end{array} \\ \left\{ \begin{array}{c} 1 & 0 \end{array} \\ 0 \end{array} \\ \left\{ \begin{array}{c} 1 & 0 \end{array} \\ 0 \end{array} \\ \left\{ \begin{array}{c} 1 & 0 \end{array} \\ 0 \end{array} \\ \left\{ \begin{array}{c} 1 & 0 \end{array} \\ 0 \end{array} \\ \left\{ \begin{array}{c} 1 & 0 \end{array} \\ 0 \end{array} \\ \left\{ \begin{array}{c} 1 & 0 \end{array} \\ 0 \end{array} \\ \\ \left\{ \begin{array}{c} 1 & 0 \end{array} \\ 0 \end{array} \\ \\ \left\{ \begin{array}{c} 1 & 0 \end{array} \\ \\ \left\{ \begin{array}{c} 1 & 0 \end{array} \\ \\ \\ \left\{ \begin{array}{c} 1 & 0 \end{array} \\ \\ \left\{ \begin{array}{c} 1 & 0 \end{array} \\ \\ 0 \end{array} \\ \\ \\ \left\{ \begin{array}{c} 1 & 0 \end{array} \\ \\ \\ \\ \left\{ \begin{array}{c} 1 & 0 \end{array} \\ \\ \\ \\ \\ \\ \\ \end{array} \\ \\ \\ \\ \\ \\ \end{array} \\ \\ \\ \\ \end{array} \\ \\ \end{array} \\ \\ \\ \\ \end{array} \\ \\ \\ \end{array}$$

$$\begin{split} & A = Y(A, S) \in (z, A; Z), E(3, I; Z), F(2, A; J), F(2, A; J), A(2, I), E(1, S, I), \\ & A = P(A, S) \in (z, A; Z), E(3, I; Z), F(2, A; J), F(2, A; J), A(2, I), E(1, S, I), \\ & E(z, S; -Z), A(3, -I) \\ & So ig A is unerticle provided we keep trade of gradients and the unbox we can be an action express A as, product of elementary needed. \\ & Comple : \\ & A = \begin{pmatrix} 0 & 0 & 1 \\ 1 & 1 & 2 \end{pmatrix}, A(2, -I) \\ & A = \begin{pmatrix} 0 & 0 & 1 \\ 1 & 1 & 2 \end{pmatrix}, B(A, S) \\ & (A = I - I) \\ & (A$$

DPcop: Suppose i Row B=O, then it Row BC= O for any C Proof: BC = E Bij Cjk = 0 BProposition: If B is a square materix and its Row B=0, then B is not Proof: Suppose B is investible, then BBT = In This implies muertible. it Row of In=0 [False] So B is not invertible QED 3 Reposizion: let A be a square mettic and X invertible matrix IS XA has a zero row then A is structible Proc Suppose A is invertible $AA^{-1} = I_A$; $(XA) \cdot A^{-1} = X$ (which is invertible). But XA has a zero row, so X has a zero row: CONTRADICTION $A = \begin{pmatrix} 1 & -1 & 1 \\ 1 & 1 & 1 \\ 1 & -3 & 1 \end{pmatrix} = \begin{pmatrix} 0 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 1 & 0 \end{pmatrix} \begin{pmatrix} 0 & 0 & 0 \\ 0 & 0 & 0 \\ -1 & 0 & 0 \\ -1 & 0 & 0 \end{pmatrix} \begin{pmatrix} 0 & 0 & 0 \\ -1 & 0 & 0 \\ 0 & 0 & 0 \end{pmatrix}$ In the example. take X = E(3,2;1)E(3,1;-1)E(2,1,-1)XA has a zero Row. X invertible so A not intertible QED "START COURSE": PROPOSITIONAL LOGIC and, or, not, implies I Start with statements which can be either True (T) or Folse (F). eg: p= t is caining, q= it is cold. ATOMIC STATEMENTS. P 9 Parel9 T T T F T F Conjunction - only is True when both are true. TRUTH TABLES 2 OR Two possible senses: 2.1 Inclusive or : Lowin = Vel 2.2 Exclusive or latin = aut

$$\begin{bmatrix} \mathbb{Z} \end{bmatrix} INCLUSIVE = \mathbb{Q} = \mathbb{V} = DISSUMEDICEN$$

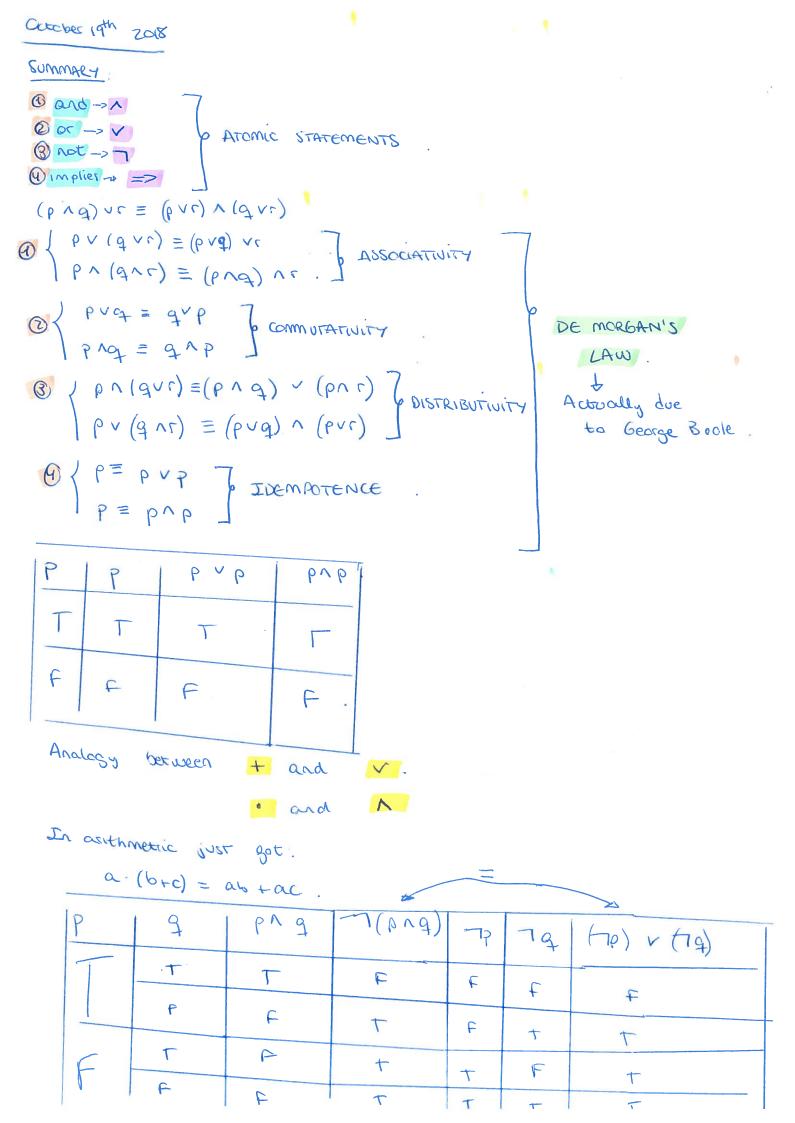
$$\frac{\mathbb{Z}}{\mathbb{Z}} = \frac{\mathbb{Z}}{\mathbb{Z}} = \mathbb{V} = \mathbb{Z} = \mathbb{$$

	bod '		(pvq) rr	pvq	(pnc) v(anc)	ρΛς	915
*	P C	7	TE		F	T F T	T F
	TF				F F	F F	F T F
	F	FFF	F	F	F	L E	F

Observe: (pVq) Nr, (pNr)V(qNr) have some truth tables TRUE (FALSE under exactly the same conductors when 2 compained statements have some Truth Table we regard them as equivalent :=

$$(p \lor q) \land r \equiv (p \land r) \lor (q \land r)$$

$$eq p = np
p T p np
T F T
F T F F
F T F$$



$$\begin{array}{c} \left(\begin{array}{c} \left(p \land q \right) = (p) \lor (p) \lor (p) \\ (p \land q) \equiv (p) \lor (p) \lor (p) \\ (p \land q) \equiv (p) \lor (p) \lor (p) \\ (p \land q) \equiv (p) \lor (p) \lor (p) \\ (p \land q) \equiv (p) \lor (p) \lor (p) \\ (p \land q) \equiv (p) \lor (p) \\ (p \land q) = (p) \lor (p) \lor (p) \\ (p \land q) = (p) \lor (p) \lor (p) \\ (p \land q) = (p) \lor (p) \lor (p) \\ (p \land q) = (p) \lor (p) \lor$$

Propositional lage ~ Contract propositions
Preduced lage ~ Lowers propositions

$$\underline{e_X} = P(X) = X \ge Z$$
.
If X is a newrow propositions $P(X)$ where X is entry 0 or 1.
So have two contracts propositions $P(X)$, $P(X)$.
 $\{0, 1\} = 20\pi ain of discussion .
1) Is $P(X)$ thus for every occurate of X³.
2) Is $P(X)$ thus for every occurate of X.
(X_X) $P(X)$ means $P(X)$ is the for every X in Domain
Existential associate
 $(X_X) P(X)$ means $P(X)$ is the for every X in Domain
 Q_X = $\{0, 1\}_X$
 $(X_X) P(X) = P(0) \land P(1)$
 $Q = \{0, 1\}_X$
 $(V_X) P(X) = P(0) \land P(1) \land P(2)$.
 $(V_X) P(X) = P(0) \land P(1) \land P(2)$.
 $(V_X) P(X) = P(0) \land P(1) \land P(2)$.
 $(V_X) P(X) = P(0) \land P(1) \land P(2)$.
 $(V_X) P(X) = P(0) \land P(1) \land P(2)$.
 $(V_X) P(X) = P(0) \land P(1) \land P(2)$.
 $(V_X) P(X) = P(0) \land P(1) \land P(2)$.
 $(V_X) P(X) = P(0) \lor P(1)$.
 $Q = \{0, 1\}_X$
 $Q = \{0, 1\}_Y$
 $Q = [N]$
 $(I_X) P(X) = P(0) \lor P(1) \lor P(2)$.
 $Q = [N]$
 $(I_X) P(X) = P(0) \lor P(1) \lor P(2)$.$

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e 💌

How does not interact with 4?

$$\begin{split} & \mathcal{A} = \begin{cases} \circ_{i} \iota_{i}^{i} & \text{at ceast } l \times n^{\alpha} \\ (\forall_{k}) \rho(k) = \rho(\alpha) \wedge \rho(4) & (\int \rho(\alpha)) \vee (\neg \rho(4)) = (\exists_{k}) (\neg \rho(k)) \\ & \mathcal{A} = \langle \sigma_{i}, z \rangle^{\alpha} \\ \neg (\forall_{k}) \rho(k) = \neg (\rho(\alpha) \wedge \rho(i) \wedge \rho(z) = \neg (\rho(\alpha) \vee \rho(i)) \vee (\neg \rho(i)) = (\exists_{k}) (\neg \rho(k)) \\ & \mathcal{A} = \langle \sigma_{i}, z \rangle^{\alpha} \\ \neg (\forall_{k}) \rho(k) = (\neg (\rho(\alpha) \wedge \rho(i) \wedge \rho(z)) = \neg (\rho(\alpha) \vee \rho(i)) \vee (\neg \rho(z)) = (\exists_{k}) (\neg \rho(k)) \\ & \mathcal{A} = \langle \sigma_{i}, \iota_{i} \rangle \\ & \mathcal{A} = \langle \rho_{i}, \iota_{i} \rangle \\ & \mathcal{A} = \langle \sigma_{i}, \iota_{i} \rangle \\ &$$

[3

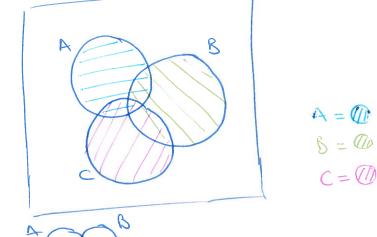
so
$$(\forall x) (\exists y) P(x, y)$$
 is not equivalent to $(\exists y) (\forall x) P(x, y)$
code of approximation in the important.
In got $(\exists y) (\forall x) P(x, y) \Rightarrow \forall x) (\exists y) P(x, y) \times$
' but converse is guide.
However,
 $(\forall x) (\forall y) = (\forall y) \forall x)$
 $(\exists x) (\exists y) = (\exists y) (\exists x)$
KIT set there one out construct
($(x) (\forall x) = (\forall y) (\forall x)$
 $(\exists x) (\exists y) = (\exists y) (\exists x)$
KIT set there one is non-backies ()
 $f_{0,1}(z) = \{1, z, 0\} = \{1, 0, 2\}$
KIT set there one out construct
() To notice also use the top backies ()
 $f_{0,1}(z) = \{4, 2\}$
o for there y has one Primitive Sight $\mathcal{E} = biologing = is a number of.$
 $X = \{0, 1, 3\}$
 $e \in X, i \in X, 3 \in X, 2 \notin X.$
 $A = for is described by its members.$
 $A = for is described by its members.$
 $A = for with f and end with f .
Stars begin with f and end with f .
Thus ways of described a last:
 $(\mathbb{C} = Matrix with f) and end with f .
 \mathbb{C} water is described f is $f = (x \in X = Y, Z \notin X)$.
 $(\mathbb{C} = \{0, 1, 2\})$
 $(\mathbb{C} = \{1, 2, 3\})$
 $(\mathbb{C} = \{1, 3\})$
 $(\mathbb{$$$

Supex A, B are sets where
$$A \subset B'$$
 when for each $X \in A$, $X \in B$
 $(\forall x) (x \in A \rightarrow x \in B)$.
Mustrates A
 \Rightarrow Don't confine membership with inclusion.
 $X = \{0, 1, \{0, 14\}, \{1, 24\}, \{14\}, \{0, 1, 24\}\}$
 $o \in X, \vee$ $\{0, 1', \{1, 24\}, \{14\}, \{0, 1, 24\}\}$
 $o \in X, \vee$ $\{0, 1', \{1, 24\}, \{14\}, \{0, 1, 24\}\}$
 $o \in X, \vee$ $\{0, 1', \{1, 24\}, \{14\}, \{0, 1, 24\}\}$
 $o \in X, \vee$ $\{0, 1', \{1, 24\}, \{14\}, \{0, 1, 24\}\}$
 $o \in X, \vee$ $\{0, 1', \{1, 24\}, \{14\}, \{0, 1, 24\}\}$
 $i \in X$ \vee $\{1, 24] \subset X'$ $i \in n$ a busch inclusion in the set = 100
 $\{14] \in X \vee$ $\{1, 24] \subset X'$ $i \in n$ a member i when i a set i by i a different gale group goes on i and i for i and i

PA(X) is the predicate which defines A.

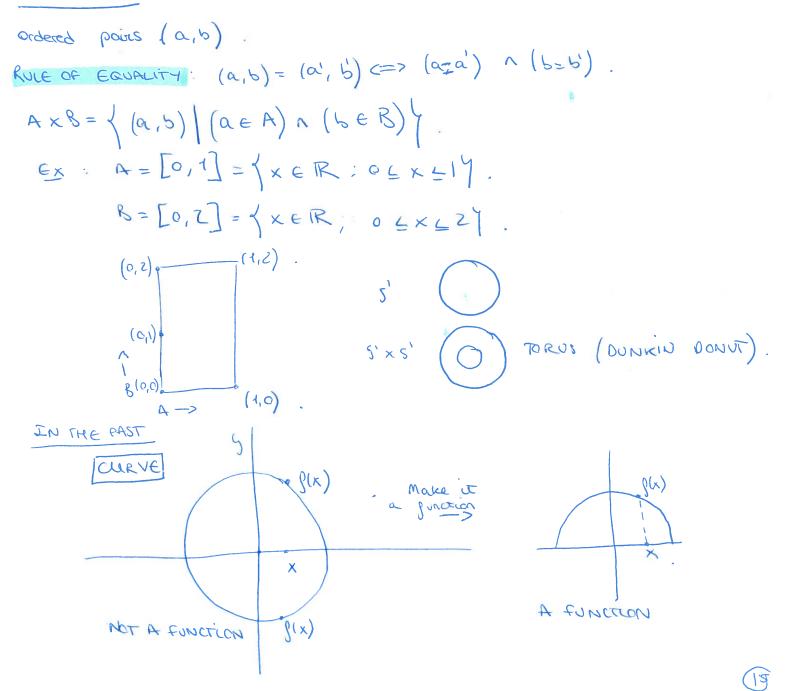
A-A = Ø





$$A \cap B =$$

PRODUCT SETS:



Lytimes dep A of forces
A is sets
By a forces
$$p: A \rightarrow B$$
.
We near a role when assign to each a large well define
cleaners $p(x) \in B$.
A is domain of f .
B is ordenain of f .
 $f(x) = 2x+1$.
 $g(x) = \frac{1}{x+1}$
 $h(x) = x^{2}$.
 $f(x) = \sqrt{x}$.
 $f(x) = x + 1$.
 $f(x) = x + 1$.
 $f(x) = x^{2}$.
 $f(x) = \sqrt{x}$.
 $f(x) = \sqrt{x}$.
 $f(x) = x^{2}$.
 $f(x) = \frac{1}{x^{2}x^{2}}$.
 $g(x) = \frac{1}{x^{2}x^{2}}$.
 $g(x) = \frac{1}{x^{2}x^{2}}$.
 $f(x) = \sqrt{-2}f - 2$.
 $f(x) = \sqrt{-2}f - 2$.
 $f(x) = x^{2}$.
 $f(x) = \sqrt{-2}f - 2$.
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 $f(x) = x^{2}$.
 $f(x) = \sqrt{-2}f - 2$.
 $f(x) = \sqrt{-2}f$

$$h(\omega) = \sqrt{x}$$
Put $R \ge 0 = \left\{ x \in \mathbb{R} : x \ge 0 \right\}$

$$h: R_{\ge 0} \longrightarrow R h(x) = \sqrt{x}$$

$$Iy we notice a convection that $\sqrt{x} \ge 0$ then we have a mapping.
$$A \stackrel{f}{\Longrightarrow} 6 \stackrel{f}{\Longrightarrow} C$$

$$I gate a mapping $g \circ g \cdot g$

$$g \circ g : A \longrightarrow C$$

$$\left[(g \circ g) (\alpha) - g (g(\alpha)) \right].$$
Convection a Awayan Associative
$$C \stackrel{h}{\longrightarrow} D$$

$$A \stackrel{f}{\longrightarrow} 8 \stackrel{f}{\longrightarrow} C \stackrel{h}{\longrightarrow} D$$

$$A \stackrel{f}{\longrightarrow} 8 \stackrel{h \circ g}{\longrightarrow} C$$

$$\left[h \circ g \circ g \right] (\alpha) = h (g \circ g(\alpha)) = h (g(g(\alpha))) = 1$$

$$\left[(h \circ g \circ g) \right] (\alpha) = (h \circ g) (g(\alpha)) = h (g(g(\alpha))) = 1$$

$$\left[(h \circ g \circ g) = (h \circ g) \circ g \cdot Q \in D \right].$$
How about emmunicativity?
$$f: R \longrightarrow R \quad g(x) = Ces(x).$$

$$\left(g \circ g (x) = (cos(x))^{2} + 1 \right] \neq \text{ They are write some}.$$

$$\left(f \circ g \circ g (x) = (cos(x))^{2} + 1 \right].$$$$$$

$(3 \circ 9)(0) = 051 < 1$. $(3 \circ 9)(0) = 2$. NOT THE SAME.	
Proposition composition is not computative usually.	
Jentity mapping If A is a set define IdA: A -> A	х. Х
$Id_A(\alpha) = \alpha$ $ex: g(x) = x$ Invertible mapping	8
Let S: A -> B be mapping. Say that g is invertible when	
there exists a mapping 9: B -> A such that P A SB -> A	
$\begin{cases} g \circ g = Jd_A; A \longrightarrow A \\ J \circ g = Jd_B; B \longrightarrow B \end{cases}$	*
Example $R_{>0} = \int x \in R$; x>04.	
$exp. \mathcal{R} \longrightarrow \mathcal{R} > 0 exp(x) = \underbrace{\sum_{r=0}^{\infty} \frac{x^r}{r!}}_{r!}$ $exp. \mathcal{R} \longrightarrow \mathcal{R} exp(x) = \int_{1}^{x} \frac{dt}{t}$	
$\begin{cases} exp \circ log(x) = x \\ log(e^{x}) = x \end{cases} \begin{pmatrix} e^{logx} = y \\ log(e^{x}) = x \end{pmatrix}.$	2
$\begin{array}{l} f: \mathcal{R} \longrightarrow \mathcal{R} f(x) = 2x + 1 g(x) = \frac{x - 1}{2} \\ \text{want } g: \mathcal{R} \longrightarrow \mathcal{R} \left(\int \circ g \right)(x) = x \\ \left(\int \circ g \right)(x) = x \end{array}$	
$\begin{cases} 0, 1, 2 \\ 1 \\ 0, 1, 2 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \\$	
gof = Id $po - > t - > 0$. log = Id.	

Find: If
$$j: A \rightarrow 8$$

has an invester mapping then the inverse mapping is unique.
Proof: Suppose $g: B \rightarrow A$ j he sufficients that does are note than
 $h: B \rightarrow A$ j hereits that does are note than
 $g \circ g = Idg$
 $h \circ j = Idg$
 $h \circ i = i dg$
 $h \circ i = i dg$

· SURSECTIVITY

A

J: A->B is subjective UbeB JaeA Ja)=b In English, everything in B can be hit: by something in A. Example :

J:
$$R \rightarrow R$$

J(x) = 7 x+1.
J: is surjective why?.
Take $y \in R$. Got to grid $x \in R$ $J(x) = J$. Take $x = (J-1)$
Example
 $g: R \rightarrow R$
 $g(x) = x^2$.
J is not surjective
 $h(z) = z^2$ then h is surjective (parque ($h(z)$ si que existe al ver
 $h(z) = z^2$ then h is surjective ($parque (h(z)$ si que existe al ver
 $h(z) = z^2$ then h is surjective ($parque (h(z)$ si que existe al ver
 $h(z) = z^2$ then h is surjective ($parque (h(z)$ si que existe al ver
 $h(z) = z^2$ then h is surjective ($parque (h(z))$ si que existe al ver
 $g: \mathbb{Z} \rightarrow \mathbb{Z}$ $g(x) = 2x^3 - x = \int_{-\infty}^{\infty} \int_$

Suppose
$$f(\alpha) = f(\alpha^{1})$$
.
Apply S. S $(f(\alpha)) = g(f(\alpha))$
 α α α .
 $S = \alpha^{1}$.
 $\begin{bmatrix} J \text{ is surjecture} \\ (at b \in B \text{ put } \alpha = g(b). \\ Apply J. $f(\alpha) = f(g(b))^{-b}$
 $f(\alpha^{1}) = b$.
 $Q(D)$.
 $Ue^{(1)}$ show the converse Every bijective mapping is invertible
FormAL DEP or inperiod the's using the formal def to prove thes
 $J: R \to R \quad f(n) = x^{2}$
 $\Sigma \quad can \quad drows as graph:
 $Uman g = R \to R$.
The graph is a subset of domain x codomain. It consists of the set.
 $f(x, f(x)) : x \in Domain Y$.
 $Def^{1}: (at A, G = be sets.)$
 $S = a mapping f: A \to B$.
 $\Sigma = mean = subset f C = A \times B$.
When consider is to following 2 containers.
 $(IJ = (\alpha, b) \in J = g(\alpha))$
 $(II) = Va = R = I = b \in B \quad (\alpha, s) \in J = p(b) = f(a) is defined.)$
 $(III) = I = (\alpha, b) \in J = \alpha and (\alpha, b) \in J = p = \alpha = \alpha$.$$

g is sorgerive when :
(III)
$$\forall b \in B \exists a \in A$$
 $(a, b) \in g \rightarrow bdas las y & preder obtained
Suppose $g < A \times S$ (on one $x \in q \in pertense
 $b : (b, a) \in g' = 2$ $(a, b) \in g$.
(CONDITIONS FOR F TO BE A BUSCOME MAPPINE:
(III) $[(a, b) \in g] \wedge [(a, b) \in g] \Rightarrow b = b$.
(III) $[(a, b) \in g] \wedge [(a, b) \in g] \Rightarrow a = a^{t}$.
(III) $[(a, b) \in g] \wedge [(a, b) \in g] \Rightarrow a = a^{t}$.
(III) $[(a, b) \in g] \wedge [(a, b) \in g] \Rightarrow a = a^{t}$.
(III) $[(b, a) \in g'] \wedge [(b, a) \in g'] \Rightarrow a = a^{t}$.
(III) $[(b, a) \in g'] \wedge [(b, a) \in g'] \Rightarrow a = a^{t}$.
(III) $[(b, a) \in g'] \wedge [(b, a') \in g'] \Rightarrow a = a^{t}$.
(III) $[(b, a) \in g'] \wedge [(b, a') \in g'] \Rightarrow a = a^{t}$.
(III) $[(b, a) \in g'] \wedge [(b, a') \in g'] \Rightarrow a = a^{t}$.
(III) $[(b, a) \in g'] \wedge [(b, a') \in g'] \Rightarrow a = a^{t}$.
(III) $[(b, a) \in g'] \wedge [(b, a') \in g'] \Rightarrow a = a^{t}$.
(III) $[(b, a) \in g'] \wedge [(b, a') \in g'] \Rightarrow a = a^{t}$.
(III) $[(b, a) \in g'] \wedge [(b', a') \in g'] \Rightarrow b = b^{t}$.
(III) $[(b, a) \in g'] \wedge [(b', a') \in g'] \Rightarrow b = b^{t}$.
(III) $[(b, a) \in g'] \wedge [(b', a') \in g'] \Rightarrow b = b^{t}$.
(IV)! $\forall a \in A \exists b \in B (b, a) \in \beta^{t}$.
(IV)! $\forall a \in A \exists b \in B (b, a) \in \beta^{t}$.
(IV)! $\forall a \in A \exists b \in B (b, a) \in \beta^{t}$.
(IV)! $\forall a \in A \exists b \in B (b, a) \in \beta^{t}$.
(IV)! $\forall a \in A \exists b \in B (b, a) \in \beta^{t}$.
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(IV)! $\forall a \in A \in A \in B^{t}$.
(IV)! $\forall a \in$$$

¥.....

Two sets A, B have the same CARDINALITY when I bijecture mapping.

· Permutations

Let n ≥ 1 be an integer and consider the set of all bijective mappings.

$$I \longrightarrow I$$

$$I \longrightarrow$$

$$O_n = \{g: \{1, \dots, n\} \longrightarrow \{1, \dots, n\}$$

$$\int_{z=1}^{z} \{J, \dots, n\} \longrightarrow \{1, \dots, n\}$$

$$\int_{z=1}^{z} \{Jd, T\} \text{ where } T(1) = 2 \quad T(z) = 1$$

$$\int_{z=1}^{z} T = 1$$

Note that TOT= Id. Proposition Let a 1, B 3, C.

be invertible mappings. Then god; A-> c is also invertible

$$\frac{(9 \circ 8)^{-1} = 8^{-1} \circ 9^{-1}}{8 \circ 8^{-7} \otimes 2^{-1} \circ 8^{-1}} \circ 8^{-1} = 9 \circ 12 \circ 8^{-1}$$

$$= 9 \circ 8^{-1} = 12$$

$$\frac{8^{-1} \circ 8^{-1} \circ 9 \circ 8}{8^{-1} \circ 8^{-1} \circ 8} \circ 8 = 8^{-1} \circ 12 \circ 8$$

$$= 8^{-1} \circ 8$$

$$= 8^{-1} \circ 8$$

$$= 5 \circ 8$$

$$= 5 \circ 8$$

$$\begin{aligned}
 \int_{3}^{2} \cos 6 & \text{elements} = 3! = 3 \cdot 2 \cdot 1 = 6 \\
 I = 1 \\
 I = -2 \\
 J = -2 \\$$

TOO(1) = 17 T O To

$$\sigma(z) = 3 \quad T \circ \sigma = \sigma \circ \tau$$

$$\sigma(3) = z \quad [\neq \sigma \circ \tau]$$

$$S_{3} = \begin{cases} \langle i \rangle, \sigma, \sigma^{2}, \tau, \sigma\tau, \sigma^{2}\tau. \rangle$$

$$S_{3}^{2} = i \quad \tau^{2} = i$$

$$\tau \sigma = \sigma^{2}\tau$$

$$(\sigma^{2} november 2018)$$

$$\frac{\delta^{2}}{\delta \sigma} + seet \times is juite when criver i) X = \emptyset \quad er$$

$$iii) = b_{jecture mapping } g: \{1, ..., n\} \to X$$

$$(s\sigma = a_{1}so \quad g^{-1} : X \to \{1, ..., n\}' : B_{ijecture}$$

$$IS \times is gritte \quad p: \{1, ..., n\}' \to B_{ijecture}$$

$$IX| = n$$

$$IX| = n$$

$$IX| = conctinue \quad g \times A$$

$$[y| = 0$$
Represention: let X, J be gritte serve with $|x| = |y| = n \ge 1$.

$$(er \quad g: X \to J \quad be \quad a_{1}nopping$$

$$Then, \quad L) \quad g \quad injecture \quad \Longrightarrow g \quad surjecture.$$

Coursely true on
$$n=1 - 2 n=1 \ge 1$$

Suppose for $n-1$.
Suppose for $n-1$.
Suppose for $n-1$.
Cant $\mathcal{R} = \{x_1, \dots, x_n\}$.
 $f(x_n) \in \mathcal{A}$.
 $\mathcal{R}_{i} = x - \{x_n \in \mathcal{A}_{i}\}$ $f(x_n) \in \mathcal{A}$.
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 $\mathcal{R}_{i} = x - \{x_n \in \mathcal{A}_{$

50 8: x -> J also surjecture.

$$\begin{array}{c} x = y_{1} & y_{2} & z_{2} & z_$$

$$J = \begin{pmatrix} 1 & 2 & \dots & n \\ p(1) & f(2) & \dots & p(N) \end{pmatrix} \quad \text{Fore From op } f.$$
There is an abbreviourbal from Explorements
$$CYCOLC \text{ Remomentants}$$

$$C_{X} : \begin{pmatrix} 1 & 2 & 3 & 4 & 5 & 6 \\ 4 & 2 & 6 & 3 & 5 & 1 \end{pmatrix}$$

$$= \begin{pmatrix} 1, 4, 3, 6 \end{pmatrix} \dots$$

$$C_{X} : \begin{pmatrix} 1 & 2 & 3 & 4 & 5 & 6 \\ 5 & C & 1 & 2 & 3 & 7 & 4 \end{pmatrix}$$

$$= \begin{pmatrix} (2, 6, 7, 4) (15, 3) \\ 5 & C & 1 & 2 & 3 & 7 & 4 \end{pmatrix}$$

$$= \begin{pmatrix} (2, 6, 7, 4) (15, 3) \\ 5 & C & 1 & 2 & 3 & 7 & 4 \end{pmatrix}$$

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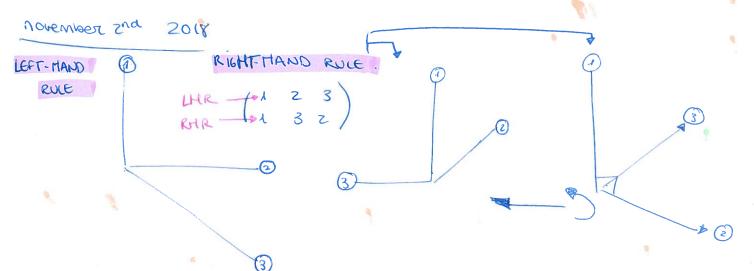
$$= \begin{pmatrix} 1 & 2 & 3 & 4 & 5 & 6 & 7 \\ 5 & C & 1 & 2 & 3 & 7 & 4 \end{pmatrix}$$

$$= \begin{pmatrix} 2, 6, 7, 4 \\ 1, 5 & 3 & 7 & 4 \end{pmatrix} (2, 5, 3) (2, 6, 7, 4) \dots$$

$$= \begin{pmatrix} 1 & 2 & 3 & 4 & 5 & 6 & 7 \\ 6 & 1 & 2 & 3 & 7 & 7 & 7 \\ 1 & 1 & 2 & 2 & 3 & 7 & 7 \\ 1 & 2 & 2 & 2 & 3 & 7 & 7 \\ 1 & 2 & 2 & 2 & 3 & 7 & 7 \\ 1 & 2 & 3 & 4 & 5 & 6 & 2 & 7 & 9 & 10 & 1 \\ 1 & 2 & 3 & 4 & 5 & 6 & 2 & 7 & 9 & 10 & 1 \\ 1 & 2 & 3 & 4 & 5 & 6 & 2 & 7 & 9 & 10 & 1 \\ 1 & 2 & 3 & 4 & 5 & 6 & 2 & 7 & 9 & 10 & 1 \\ 1 & 2 & 3 & 4 & 5 & 6 & 2 & 7 & 9 & 10 & 1 \\ 1 & 2 & 3 & 4 & 5 & 6 & 2 & 7 & 9 & 10 & 1 \\ 1 & 2 & 3 & 4 & 5 & 6 & 2 & 7 & 9 & 10 & 1 \\ 1 & 2 & 3 & 4 & 5 & 6 & 2 & 7 & 9 & 10 & 1 \\ 1 & 2 & 3 & 4 & 5 & 6 & 2 & 7 & 9 & 10 & 1 \\ 1 & 2 & 2 & 4 & 5 & 6 & 2 & 7 & 9 & 10 & 1 \\ 1 & 2 & 3 & 4 & 5 & 6 & 2 & 7 & 9 & 10 & 1 \\ 1 & 2 & 3 & 4 & 5 & 6 & 2 & 7 & 9 & 10 & 1 \\ 1 & 2 & 2 & 4 & 5 & 6 & 2 & 7 & 9 & 10 & 1 \\ 1 & 3 & 7 & 7 & 1 & 1 & 1 & 0 & 5 & 6 & 4 & 2 & 9 \end{pmatrix}$$

.

$$= (4,11,9)(1,3,7,5)(2,8,6,10) = (1,3,7,5)(2,8,6,10)(4,11,9)$$



To each permutation.

$$g = \begin{pmatrix} 1 & z & \ddots & \ddots \\ g(1) & g(2) & \cdots & g(n) \end{pmatrix}$$

we associate a sign +1 or -1

$$H \longrightarrow preserves oriensation
-1 \longrightarrow reverses arentation
Example $G = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 \\ 2 & 3 & 4 & 5 \\ 2 & 1 & 3 & 4 & 5 \\ 2 & 1 & 3 & 4 & 5 \\ 2 & 1 & 3 & 4 & 5 \\ 2 & 3 & 1 & 4 & 5 \\ 2 & 3 & 1 & 4 & 5 \\ 2 & 3 & 1 & 4 & 5 \\ 2 & 3 & 4 & 1 & 5 \\ 2 & 3 & 4 & 1 & 5 \\ 2 & 3 & 4 & 1 & 5 \\ 2 & 3 & 4 & 1 & 5 \\ 2 & 3 & 4 & 1 & 5 \\ 2 & 3 & 4 & 1 & 5 \\ 2 & 3 & 4 & 1 & 5 \\ 2 & 3 & 4 & 1 & 5 \\ 2 & 3 & 4 & 1 & 5 \\ 2 & 3 & 4 & 1 & 5 \\ 2 & 3 & 4 & 5 \\ 2 & 3 & 4 & 5 \\ 2 & 3 & 4 & 5 \\ 2 & 3 & 4 & 5 \\ 2 & 3 & 4 & 5 \\ 2 & 5 & 5 & 5 & 5 & 5 \\ 2 & 3 & 4 & 1 & 5 \\ 2 & 5 & 5 & 5 & 5 & 5 & 5 \\ 2 & 3 & 4 & 1 & 5 \\ 2 & 5 & 5 & 5 & 5 & 5 & 5 \\ 2 & 3 & 4 & 1 & 5 \\ 2 & 5 & 5 & 5 & 5 & 5 & 5 \\ 2 & 3 & 4 & 1 & 5 \\ 2 & 5 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 & 5 \\ 3 & 5 & 5 & 5 &$$$

et, : Generalisation $(a_1, a_2, a_3, ..., a_n) = (a_1, a_n) (a_1, a_{n-1}) ... (a_1, a_3)(a_1, a_2)$ $c_{ycle of length n = (n-1)$ there are (n-1) multiplication terms. A transpisution is a cycle of length 2.

ex: (z, 7) o (7, 100). Proposition A cycle of length is a product of (n-1) transpositions

50:
3) A cycle of all length is a prawe of any Even no of transportant
1) A cycle of even length is a prawe of any odd no of transportant
Last time we showed that:
Therein Any permutation is a prawe of object cycles.
Contact: Any permutation is a prawe of transportant
Any permutation of the prawe of the prawe of even and also
as prawe of all no prawe of any order of the prawe of the second
of
$$(A = 2 - 1)$$
 when O is prawe of any of the order to write
Some permutation as been preduce of all the prawe to write
Some permutations
Mean to assume the frant is interpret prawe to write
Some permutation of the prawe of a permutation
 $A = (A = 2 - 1) + S = A = 2 + (O + 1)$
 $A = (A = 2 - 3 + 5 = A + S = (A = 0 + 1) + (A = 1) + (A$

.

ORDER OF PERMUTATION ord (σ) = mim $\{N \ge 1\}$; $\sigma^{N} = 1$ example: (1,5,3,4) has order 4. Grover of (a, az, ... an) = n . Iterate a cycle of length n, n times and get the identity (1,5,3,4). Volver a consegur el = nº del principio example: order 4 order 3 G = (i, 5, 3, 4) (2, 7, 6)X, Xz X, Xz X, Z = X2(A) (Disjoint cycles) acdec 12 occder 12 - > porque el mcm de 4y 3 es 12. Proposition. IJ X1, X2 are disjoint cycles $(x_{1}, x_{2})^{N} = X_{1}^{N} X_{2}^{N}$. (This is only true because) 1 1 $X_2 X_1 = X_1 X_2$ Ropisition: IJ X1, ..., Xm are disjoint cycles multiple. Order (X, Xm) = (LCM (ord(X)), ..., ord (Xm)) = LCM (length(x,), ... length (xm)), (X, Xm) = X, X2 Xm Because they commute. $\sigma = (x, 5, 8, 11, 9)(2, 7, 3)(4, 10, 12, 6) \quad sign \Theta = -1$ 0(d(e)) = LCM(5,3,4) = 60Say that a transposition (1, j) is ADJACENT when [-i]=1 e_X : (3,4) (6,7). but not (1,3) . Proposition: Any transposition is a product of an ODO no of adjacent 1 1 transportance Roog: Degine gap (i,j) = 1j-i) gap (i, j) = k then (i, j) is a product of (2k-1) adjacent

transpositions

$$\begin{aligned} \begin{split} & = \begin{cases} 0, 1, 2, \dots, n+1, \dots, 1 & \text{is a nepping of a product of a constrained of a const$$

$$\begin{aligned} \forall x \in IF \quad \exists (-x) \in F \quad x + (-x) = (-x) + x = 0 \quad \text{ADDITIVE INVERSE} \\ \text{MOUTIPULCATIVE AXIENS:} \\ \bullet, \quad F \times F \quad \longrightarrow IF \\ (x, y) \quad \longrightarrow x \cdot y \quad (\text{Act} \quad (K, y)) \\ \text{Such that} \quad (x \cdot y) \cdot z = \quad x \cdot (y \cdot z) \quad \text{Associative} \\ & x \cdot y = y \cdot x \quad \text{Computerative} \\ & x \cdot y = y \cdot x \quad \text{Computerative for with} \\ & x \cdot 1 = 1 \cdot x = x \quad 1 \quad \text{is} \quad \text{MULTIPLICATIVE for with} \\ & \forall x \in IF + foy \exists \ x^{-1}; \quad X \cdot x^{-1} = x^{-1} \cdot x = 1 \quad \text{MULTIPULATIVE} \\ & (\text{Non} \quad \text{zero elements have multiplicative} \quad \text{Friends} \\ & \text{MULTIPLICATIVE} \quad \text{Axiem} \end{aligned}$$

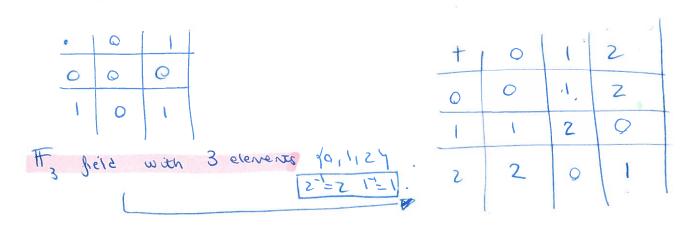
$$\mathbb{P} = (\mathbb{Q}, +, 0, \cdot, 1) \text{ is } \mathbb{P} = (\mathbb{Q}, +, 0, \cdot, 1) \text{ is } \mathbb{P} = \mathbb{P} = (\mathbb{Q}, +, 0, \cdot, 1) \text{ is } \mathbb{P} = \mathbb{P$$

my

EF

IF2 Sield with 2 elements. $f_z = \{even, 000\}, 000 = 1, prive et resultado es un n° par pargo 0,$ $<math>f_z = \{even, 000\}, 000 = 1, prive et resultado = odial pargo -1.$

+	EVEN	900		+	0)
Even	EVEN	000	-7	0	0	X
000	000	EVEN				0



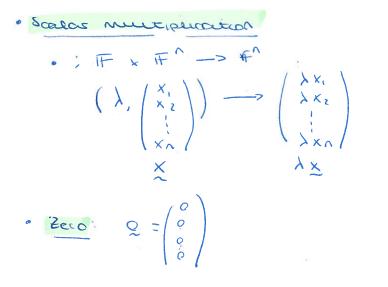
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November 14th 2018

 $TF \text{ field} \\ TF^{n} = \int X = \begin{pmatrix} X_{1} \\ X_{2} \\ \vdots \\ X_{n} \end{pmatrix}; X_{n} \in TF$

· Addition

 $\begin{array}{c} + \cdot & F^{n} \times F^{n} \rightarrow iF^{n} \\ \begin{pmatrix} x_{1} \\ x_{2} \\ \vdots \\ x_{n} \end{pmatrix} \cdot \begin{pmatrix} y_{1} \\ y_{2} \\ \vdots \\ y_{n} \end{pmatrix} \rightarrow \begin{pmatrix} x_{1} + y_{1} \\ x_{2} + y_{2} \\ \vdots \\ x_{n} + y_{n} \end{pmatrix}$



I ADDIFINE PROPERTIES:

 $^{\circ}$ $\stackrel{\times}{\rightarrow}$ + $(\stackrel{\vee}{2} + \stackrel{Z}{=}) = (\stackrel{\times}{2} + \stackrel{\vee}{2}) + \stackrel{Z}{=} : Associative$ · X + y = y + x : commutative $\circ \circ + \chi = \chi + \circ = \chi$: Identity • $\forall x \in \mathbb{F}^{n}$ $\exists -x \in \mathbb{F}^{n}$; x + (-x) = Q: inverses 2) SCALAR MULTIPLICATION PROPERTIES · $\lambda \cdot (x + z) = \lambda \cdot x + \lambda \cdot y$ distributive · (\ + m) · x = \x + mx : distributive · A X = X identicy " O · X = O · Inverse > (m x) = (hu) x assurative Degni Let IF be a field. By a vector space. Vover IF. I mean $V = (V, + 0, \cdot)$ where: i) V is a set. $ii) o \in V$ iii) +; V × V -> V is a mapping which sort isgies the properties above. (×, ×) -> ×+3 • x + (y + z) = (x + y) + z Associative · × + > = > + × · commercie · OTX = X+Q = X identity $\forall x \in V, \exists (-x) \in V : x + (-x) = 2 : nivers.$

Example:
i) IF
$$h$$
 is a veccor space over F for all $n \ge 2$
ii) For $n=1$, F is a veccor space over F .
Snag: Note every vector space is accountly on F^n . $F^d = rockarle
Example $F = 0$
 $V = \{\begin{pmatrix} \lambda \\ -\lambda \end{pmatrix}\}; \lambda \in F^d$.
 $V = \{\begin{pmatrix} \lambda \\ -\lambda \end{pmatrix}\}; \lambda \in F^d$.
 $V = \{\begin{pmatrix} \lambda \\ -\lambda \end{pmatrix}\}; \lambda \in F^d$.
 $V = \{\begin{pmatrix} \lambda \\ -\lambda \end{pmatrix}\}; \lambda \in F^d$.
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 $V = \{\begin{pmatrix} \lambda \\ -\lambda \end{pmatrix}\}; \lambda \in F^d$.
 $V = \{\begin{pmatrix} \lambda \\ -\lambda \end{pmatrix}\}; \lambda \in F^d$.
 $V = \{\begin{pmatrix} \lambda \\ -\lambda \end{pmatrix}\}; \lambda \in F^d$.
 $V = \{\begin{pmatrix} \lambda \\ -\lambda \end{pmatrix}\}; \lambda \in F^d$.
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 $V = \{\begin{pmatrix} \lambda \\ -\lambda \end{pmatrix}\}; \lambda \in F^d$.
 $V = \{\begin{pmatrix} \lambda \\ -\lambda \end{pmatrix}\}; \lambda \in F^d$.
 $V = \{\begin{pmatrix} \lambda \\ -\lambda \end{pmatrix}\}; \lambda \in F^d$.
 $V = \{\begin{pmatrix} \lambda \\ -\lambda \end{pmatrix}\}; \lambda \in F^d$.
 $V = \{\begin{pmatrix} \lambda \\ -\lambda \end{pmatrix}$$

· What the degn says is that I'v', ..., Yn Y is L. I if the only way to get a is by hering all adopticients =0.

Example V = IF 3

$$\frac{e_{11}}{e_{12}} = \begin{pmatrix} 1 \\ 0 \\ 0 \end{pmatrix} \qquad e_{12} = \begin{pmatrix} 0 \\ 1 \\ 0 \end{pmatrix} \qquad e_{13} = \begin{pmatrix} 0 \\ 0 \\ 1 \end{pmatrix} \qquad e_{13} = \begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix} \qquad e_{13} = \begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix} \qquad e_{13} = \begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix} \qquad e_{13} = \begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix} \qquad e_{13} = \begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix} \qquad e_{13} = \begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix} \qquad e_{13} =$$

claim that
$$\int e_1 \cdot e_2 \cdot e_3 \forall is L.\Sigma$$

so suppose $\cdot \cdot \lambda_1 e_1 + \lambda_2 e_2 + \lambda_3 e_3 = 0 = \begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix}$
 $\lambda_1 \begin{pmatrix} i \\ 0 \\ 0 \end{pmatrix} + \lambda_2 \begin{pmatrix} 0 \\ i \\ 0 \end{pmatrix} + \lambda_3 \begin{pmatrix} 0 \\ i \\ 0 \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix}$
 $\begin{pmatrix} \lambda_1 \\ 0 \\ 0 \end{pmatrix} + \begin{pmatrix} 0 \\ \lambda_2 \\ 0 \end{pmatrix} + \begin{pmatrix} 0 \\ \lambda_3 \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix}$
 $\begin{pmatrix} \lambda_1 \\ \lambda_2 \\ \lambda_3 \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix}$ and read $\begin{pmatrix} \lambda_1 = 0 \\ \lambda_1 = 0 \end{pmatrix}$, $\lambda_2 = 0$, λ_3

Example ? V= IF 3.

$$\begin{aligned} & \varphi_1 = \begin{pmatrix} i \\ 0 \\ 0 \end{pmatrix} & \varphi_2 = \begin{pmatrix} i \\ i \\ 0 \end{pmatrix} & \varphi_3 = \begin{pmatrix} i \\ i \\ i \end{pmatrix} \\ & \varphi_3 = \begin{pmatrix} i \\ i \\ i \end{pmatrix} \\ & \varphi_3 = \begin{pmatrix} i \\ i \\ i \end{pmatrix} \\ & \varphi_3 = \begin{pmatrix} i \\ i \\ i \end{pmatrix} \end{aligned}$$

Take

$$\lambda_{1} \begin{array}{c} \varphi_{1} + \lambda_{2} \begin{array}{c} \varphi_{2} + \lambda_{3} \end{array} \begin{array}{c} \varphi_{3} = \varphi_{1} \\ \begin{pmatrix} \lambda_{1} \\ \varphi \\ \varphi \end{array} \right) + \begin{pmatrix} \lambda_{2} \\ \lambda_{3} \\ \varphi \end{array} \right) + \begin{pmatrix} \lambda_{3} \\ \lambda_{3} \\ \varphi \end{array} \right) = \begin{pmatrix} \varphi \\ \varphi \\ \varphi \end{array} \right) \\ \begin{pmatrix} \lambda_{1} + \lambda_{2} + \lambda_{3} \\ \lambda_{2} + \lambda_{3} \\ \varphi \end{array} \right) = \begin{pmatrix} \varphi \\ \varphi \\ \varphi \end{array} \right) \\ \begin{pmatrix} \varphi \\ \varphi \\ \varphi \end{array}$$

read: $\lambda_3 = 0$ $\lambda_2 + \lambda_3 = 0$; $\lambda_2 = 0$ $\lambda_1 + \lambda_2 + \lambda_3 = 0$; $\lambda_1 = 0$. v = 0

Example:
$$V = \mathbb{F}^{3}$$

 $\mathcal{F}_{i} = \begin{pmatrix} i \\ j \end{pmatrix}$ $\mathcal{F}_{2} = \begin{pmatrix} i \\ -i \end{pmatrix}$ $\mathcal{F}_{3} = \begin{pmatrix} i \\ 2 \end{pmatrix}$
is not L.T:
 $2 + i - \frac{1}{2} = \frac{1}{2}$ so we have
 $50 = 2\frac{1}{2}i - \frac{1}{2}i - \frac{1}{2}i = 0$.
and at least one coefficient $\neq 0$
 $\lambda_{i} = 2, \lambda_{i} = -i$, $\lambda_{i} = -i$.
 $\begin{bmatrix} \lambda_{i} \neq 0 \end{bmatrix}$
Vector space / \mathbb{F}_{i} takes les vectores del espais vectorise de abtienen
Vector space / \mathbb{F}_{i} takes les vectores indep
 $\text{Ger} = \bigvee_{i, \dots, N_{n}} \in V$. Obtainer
 $\psi_{i} \in V, \exists \lambda_{i}, \dots, \lambda_{2} \in \mathbb{F}_{n}$ is $V = \lambda_{i}V_{i} + \dots + \lambda_{n}V_{n}$
In english: every vector χ is $V \cdot C$ to complet de the detrict de
 $\text{Gar be expressed as a Livere combination in $\begin{cases} V_{1}, \dots, V_{n} \\ V_{n}, \dots, V_{n} \end{cases}$.
 $\sum_{n} \in \mathbb{F}^{3} : w_{n}V^{2}$.
 $\chi \in \mathbb{F}^{3} : \chi = \begin{pmatrix} i \\ 0 \\ 0 \end{pmatrix} = e^{i} = \begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix} = e^{i} = \begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix}$.
 $\chi = \kappa_{i}e_{i} + \kappa_{2}e_{i} + \kappa_{3}e_{i}s$.
 $\widehat{F}_{\text{reccuse}} : \forall_{i} = \begin{pmatrix} i \\ 0 \\ 0 \end{pmatrix} = \mathbb{F}^{2} = \begin{pmatrix} i \\ 0 \end{pmatrix} = \mathbb{F}^{3}$.
Son bare de \mathbb{F}^{3} .$

November 16th 2018

V = veccos space /F. $V_{1}, V_{2}, \dots, V_{n} \in V.$ $\{V_{1}, \dots, V_{n} \in V.$ $\{V_{1}, \dots, V_{n} \in L.I \text{ when } : \lambda_{1} V_{1} + \lambda_{2} V_{2} + \lambda_{3} V_{3} = 0\}; (\lambda \in F).$ $[\lambda_{1} = \lambda_{2} = \lambda_{3} = 0].$ $\{V_{1}, \dots, V_{n} \in F. S.E.$ $X = \lambda_{1} V_{1} + \dots + \lambda_{n} V_{n}.$

STANDARD EXAMPLES. $V = IF^{n} = \begin{cases} x = \begin{pmatrix} x_{i} \\ x_{i} \end{pmatrix}; x_{i} \in F \\ \vdots \\ x_{n} \end{cases}$ $e_{1} = \begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix} \quad e_{2} = \begin{pmatrix} 0 \\ 1 \\ 0 \end{pmatrix} \quad e_{2} = \begin{pmatrix} 1 \\ 1 \\ 0 \end{pmatrix} \quad e_{2} = \begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix} \quad e_{2} = \begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix}$ CS ney. $\underbrace{e}_{1} = \begin{pmatrix} 1 \\ 0 \\ 0 \\ 0 \end{pmatrix} \quad \underbrace{e}_{2} = \begin{pmatrix} 0 \\ 1 \\ 0 \\ 0 \end{pmatrix} \quad \underbrace{e}_{3} = \begin{pmatrix} 0 \\ 0 \\ 1 \\ 0 \end{pmatrix} \quad \underbrace{e}_{4} = \begin{pmatrix} 0 \\ 0 \\ 0 \\ 0 \\ 0 \end{pmatrix}$ Proposition fei, ..., en jis always L.J. Proof Suppose $\lambda_1 e_1 + \cdots + \lambda_n e_n = 0 = \begin{pmatrix} 0 \\ 0 \\ \vdots \end{pmatrix}$ $\lambda_{1}e_{1} = \begin{pmatrix} \lambda_{1} \\ 0 \\ \vdots \\ 0 \end{pmatrix} \qquad \lambda_{2}e_{2}z \begin{pmatrix} 0 \\ \lambda_{2} \\ \vdots \\ 0 \end{pmatrix}$ So $\lambda_{ie_{1}} + \dots + \lambda_{ne_{n}} = \begin{pmatrix} \lambda_{i} \\ \lambda_{i} \end{pmatrix}$ Se if $\lambda_1 e_1 + \lambda_2 e_2 + \dots + \lambda_n e_n = 0 = \begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix}$ Then, reading d, = 0, dz = 0, dn=0 le le fernent is L.I

Proposition
$$\begin{cases} k_{1}^{n}, \dots, k_{n}^{n} \\ \end{cases}$$
 always spars \mathbb{F}^{n} .
Proof: $\underline{x} = \begin{pmatrix} x_{1} \\ x_{2} \\ x_{n} \end{pmatrix} = x_{1} \begin{pmatrix} 1 \\ 0 \\ 0 \\ 0 \end{pmatrix} + x_{2} \begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix} + x_{3} \begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix} + x_{4} \begin{pmatrix} 0 \\ 0$

$$\begin{aligned} \frac{1}{k_{x}} &= 0 \\ \frac{Example I}{V_{x}} & \left\{ \begin{pmatrix} x_{x} \\ x_{x} \end{pmatrix} \in 0 \\ \vdots & x_{1} + k_{2} = 0 \\ \end{bmatrix} \\ \frac{1}{V_{x}} &= 0 \\ \frac{1}{V_{$$

• Secret multiplication

$$\begin{array}{c} x = \begin{pmatrix} x_{1} \\ y_{1} \end{pmatrix}; \quad x_{1} + x_{1} + x_{2} = 0 \quad \text{feco} \quad Q = \begin{pmatrix} 0 \\ 0 \end{pmatrix} \in V \\
 & x \in V \\
\end{array}$$

$$\begin{array}{c} \lambda_{1} = \begin{pmatrix} x_{1} \\ y_{2} \end{pmatrix}; \quad x_{1} + x_{1} + x_{2} = 0 \quad \text{feco} \quad Q = \begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix} \in V \\
\end{array}$$
All other axions act scalifies as codulated already in Q³.
Reases for V

$$\begin{array}{c} y_{1} = \begin{pmatrix} -\lambda_{1} \\ y_{2} \end{bmatrix} = \begin{pmatrix} 0 \\ -\lambda_{2} \\ y_{1} \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix} \\
\end{array}$$
Reach: $\lambda_{1} = 0 \quad \lambda_{2} = 0 \\
\hline & \lambda_{2} + \lambda_{1} \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \\ 0 \\ 0 \\ \end{array}$
Reach: $\lambda_{1} = 0 \quad \lambda_{2} = 0 \\
\hline & \lambda_{2} + \lambda_{1} \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \\ 0 \\ 0 \\ \end{array}$
Reach: $\lambda_{1} = 0 \quad \lambda_{2} = 0 \\
\hline & \lambda_{2} + \lambda_{1} \end{pmatrix} = \begin{pmatrix} x_{1} \\ x_{2} \\ x_{3} \\ \end{array}$
So support X:
$$\begin{array}{c} -\lambda_{1} & y_{1} \\ \lambda_{2} + \lambda_{1} \end{pmatrix} = \begin{pmatrix} x_{1} \\ \lambda_{2} \\ \lambda_{3} \\ \end{array}$$
So support X:
$$\begin{array}{c} x_{1} \\ x_{2} \\ x_{3} \\ \end{array}$$
So support X:
$$\begin{array}{c} x_{1} \\ x_{2} \\ x_{3} \\ \end{array}$$
So support X:
$$\begin{array}{c} x_{1} \\ x_{2} \\ x_{3} \\ \end{array}$$
So support X:
$$\begin{array}{c} x_{1} \\ x_{2} \\ x_{3} \\ \end{array}$$
So support X:
$$\begin{array}{c} x_{1} \\ x_{2} \\ x_{3} \\ \end{array}$$
So support X:
$$\begin{array}{c} x_{1} \\ x_{2} \\ x_{3} \\ \end{array}$$
So support X:
$$\begin{array}{c} x_{1} \\ x_{2} \\ x_{3} \\ \end{array}$$
So support X:
$$\begin{array}{c} x_{1} \\ x_{2} \\ x_{3} \\ \end{array}$$
Want to choose $\lambda_{1}, \lambda_{2} \in \mathbb{F}$
So $\frac{1}{8} + \lambda_{2} \times 2 \\$
then $x_{2} = \Lambda_{1} + \lambda_{2} \\$

$$\begin{array}{c} x_{1} \\ x_{2} \\ x_{3} \\ \end{array}$$
Want over (F (ax_{1}))(x_{1} + x_{1} + x_{2} + x_{1} + \lambda_{1} + x_{2} + \dots + a_{1} + x_{1} = 0 \\

$$\begin{array}{c} x_{1} \\ x_{2} \\ x_{3} \\ \end{array}$$

$$\begin{array}{c} x_{1} \\ x_{2} \\ x_{3} \\ \end{array}$$
So support X:
$$\begin{array}{c} x_{1} \\ x_{2} \\ x_{3} \\ \end{array}$$

$$\begin{array}{c} x_{1} \\ x_{2} \\ x_{3} \\ \end{array}$$

$$\begin{array}{c} x_{1} \\ x_{2} \\ \end{array}$$

$$\begin{array}{c} x_{1} \\ x_{2} \\ x_{3} \\ \end{array}$$

$$\begin{array}{c} x_{1} \\ x_{2} \\ x_{3} \\ \end{array}$$

$$\begin{array}{c} x_{1} \\ x_{2} \\ x_{3} \\ \end{array}$$

$$\begin{array}{c} x_{2} \\ x_{3} \\ \end{array}$$

$$\begin{array}{c} x_{1} \\ x_{2} \\ \end{array}$$

$$\begin{array}{c} x_{1} \\ x_{2} \\ \end{array}$$

$$\begin{array}{c} x_{1} \\ x_{2} \\ x_{3} \\ \end{array}$$

$$\begin{array}{c} x_{1} \\ x_{2} \\ \end{array}$$

$$\begin{array}{c} x_{1}$$

Proposition KA is a vector space / F.

· Addition Suppose X, X E KA.

· Scales muliplication x e KA; X e FF.

$$A_{\chi} = 0$$

 $\lambda (A_{\chi}) = 0 => A(\lambda_{\chi}) = 0$
 $\sum_{\lambda} k_{\lambda} \in k_{\lambda}$

tero: AQ=Q 60 Q EKA.

Solve
$$A_{X} = 0$$
.
General solution
 $X = \begin{pmatrix} -x_{2} - x_{4} - x_{6} - x_{7} \\ x_{2} - x_{4} - x_{6} - x_{7} \\ x_{2} - x_{4} - x_{7} \\ x_{4} - x_{7} \\ x_{4} + x_{6} - x_{7} \\ x_{7} \\ x_{6} - x_{7} \\ x_{7} \\$

$$A' = \begin{pmatrix} 1 & (& 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & -1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 & 1 & 0 \\ \hline (X_1, X_2, K_3, X_4, (X_7), X_6, X_4) \end{pmatrix}$$

General solution to AX=0.

$$\begin{array}{c}
X = \begin{pmatrix}
-X_2 & -X_5 & -X_7 \\
X_2 \\
+X_5 \\
-X_5 & -X_6 \\
X_5 \\
X_6 \\
X_7
\end{array}$$

Make the dovidous choices . M7 1

$$I_{V} = X_{0}E_{1} + X_{7}E_{2} + X_{6}E_{3} + X_{7}E_{4}$$

$$Y = \begin{cases} X \in Q^{10} \\ X_{1} + X_{7}E_{2} + X_{6}E_{3} + X_{7}E_{4} \\ X_{1} + X_{7}E_{4} + X_{1}E_{7} \\ X_{1} + X_{1} \\ X_{1} + X$$

Moreover 22nd
Knows The Let V be a vector space / IF.
(V+0)
DThere exists at least 4 basis gx V.
2) Any Z bases for V have are same no of elements
dim
$$F(V) = nb of such elements$$
.
Receive easily December.
Received Let (E, ..., En Y be a basis for V.
SJ X e V, then X has a unique expression as a linear embracien
X = X, E, +... + Xn En
Received X = X(E, +... + Xn En for some Xi.
Chain this expression is only a
 $X = X, E_1 + ... + Xn En for some Xi.$
Chain this expression is only a
 $X = X, E_1 + ... + Xn En for some Xi.$
Chain this expression is only a
 $X = X, E_1 + ... + Xn En for some Xi.$
Chain this expression is only a
 $X = X, E_1 + ... + (Xn - Xn) En$
 $X = X = (X_1' - X_1) E_1 + ... + (Xn' - Xn) En$
 $X = X = (X_1' - X_1) E_1 + ... + (Xn' - Xn) En$
 $X_1' - X_1 = 0$... $Xn' - Xn = 0$.
So now
 $(X_1' - X_1) E_1 + ... + (Xn' - Xn) En = 2$.
Sur $E_1, ..., En are LE. For
 $X_1' - X_1 = 0$... $Xn' - Xn = 0$.
 $E_1 = {b \choose 0} e_2 = {b \choose 0} e_3 = {b \choose 1} = 0$...
 $F_1 = {b \choose 1} e_2 = {b \choose 1} = F_3 = {b \choose 1}$... F_3 .$

$$\begin{split} \lambda_{-} &= \begin{pmatrix} \lambda_{-} \\ \chi_{-} \\$$

Standard example

Let
$$A = (a,j)_{A \leq i \leq m}$$
 be max maxim /F
I $\leq j \leq n$ for the maxim /F
(answere T_A ; $F = F = F = F$
defined by $T_A(x) = A_X$, $F = A_X = F = (a_{1,X_1} + a_{1,2}X_1 + \dots + a_{n,X_n})$
i.e.:
 $T_A(x) = (a_{1,1} \dots + a_{n,n}) + (x_1)_{X_n} = (a_{1,X_1} + a_{1,2}X_1 + \dots + a_{n,X_n})_{X_n} + (x_1 + x_1 + x_1 + x_1 + x_1)_{X_n}$
Then T_A is a lineer
Using?:
 $T_A(X) = A(X) = A(X + y) = A_X + A_Y = T_A(A) + T_A(Y)_{X_1}$
Then T_A is a lineer
Using?:
 $T_A(X) = A(XA) = A = h T_A(X)_{X_1}$
 $Difference a great.$
 $Put P_A(F) = \{a_0(1 + a_1, h + a_2, X + \dots + a_n, A^n, Y) where $a_1 \in F$.
 $U(3, 2, 1)$ by $X = 3 \text{ sind} f$
 $A_X = b T_A(X)_{X_1}$
 $A_X = b T_A(X)_{X_1}$
 $A_X = b T_A(X)_{X_1}$
 $A_X = b T_A(X)_{X_1}$
 $Put P_A(F) = power a great.$
 $A_X = b T_A(X)_{X_1}$
 $A_X =$$

Agne D.
$$P_{n}(iF) \longrightarrow P_{n}(iF)$$

 $D (Q(n)) = (\frac{dn}{dn})$
 $e^{iF} (1) = 0$
 $D(x^{2}) = 1$
 $D(x^{2}) = 2.5 = .$
Then
 $D(x^{2}) = \frac{d}{dx} (a+b) = \frac{da}{dx} + \frac{db}{dx} = D(a) + D(b)$
 $D(\lambda = \lambda D(a) (\lambda \in F)$
The matrix of a lower matrix
 $T - V \longrightarrow W$ events
basis $\mathcal{E} = \{F_{11}, \dots, F_{n}\}$ basis for V
 $\overline{W} = \{9_{11}, \dots, 9_{m}\}$ basis for W
Consider $T(E_{1}) \in W$
 $f = \{9_{11}, \dots, 9_{m}\}$ basis for W
Consider $T(E_{1}) \in W$
 $f = \{9_{11}, \dots, 9_{m}\}$ basis for W
Consider $T(E_{1}) = W$
 $f = \{9_{12}, \dots, 9_{m}\}$ basis for W
 $F = T(E_{1})$ has a unique expression as a lower containation $m\{F_{11}, 9_{12}\}$
 $F = T(E_{1})$ basis $F = W$
 $f = \{1, \dots, 2^{m}\}$ basis for W
 $F = T(E_{1})$ basis $F = W$
 $F = \{1, \dots, 2^{m}\}$ basis $F = W$
 $F = T(E_{1}) = 2 P_{1} + 2P_{2} + \dots + 2P_{2} P_{2} - 2P_{1} P_{2}$
How do E norme conforming?
Need 2 indices
 $T(E_{1}) = Q_{11} e^{P_{1}} + Q_{22} e^{P_{1}} + \dots + Q_{m} P_{m}$
 $T(E_{1}) = Q_{12} e^{P_{1}} + Q_{23} e^{P_{2}} + \dots + Q_{m} P_{m}$
 $T(E_{1}) = Q_{12} e^{P_{1}} + Q_{23} e^{P_{2}} + \dots + P_{m} a_{m1} - P_{1}$ convertion
 $T(E_{1}) = \sum_{i=1}^{m} a_{ij} e^{P_{1}}$
 $e^{P_{1}}$ reverse E is a factorized in the factori

Def
$$\mathcal{E} = \begin{cases} \mathcal{E}_{1}, \dots, \mathcal{E}_{N} \\ \Psi = \{ \mathcal{P}_{1}, \dots, \mathcal{P}_{N} \} \text{ basis for } \mathcal{W} \\ \mathcal{T} = \begin{cases} \mathcal{E}_{1}, \dots, \mathcal{P}_{N} \\ \mathbb{I} \end{cases} \text{ basis for } \mathcal{W} \\ \mathcal{T} = \begin{cases} \mathcal{E}_{1}, \dots, \mathcal{P}_{N} \\ \mathbb{I} \end{cases} \text{ basis for } \mathcal{W} \\ \mathcal{T} = \begin{cases} \mathcal{E}_{1} \\ \mathcal{E}_{2} \end{cases} \\ \mathcal{E}$$

$$E_{S} : \Gamma(1, 1), \Gamma(1, 2), \Gamma(1, 3)$$

$$\Gamma(2, 1), \Gamma(2, 2), \Gamma(2, 3)$$

$$\sum_{j=1}^{3} \Gamma(1, j) = (\Gamma(1, 1) + \Gamma(1, 2) + \Gamma(1, 3)) + (\Gamma(2, 1) + \Gamma(2, 2))$$

$$= \Gamma(2, 3)$$

$$F(2, 3) + \Gamma(2, 3)$$
Decentionsion of a threas mapon a vasis
$$T: V \longrightarrow W \quad \text{lineas}$$
Let $Y = K_{1}, \dots, K_{n} Y$ be a vasis for V
Reposition T is completely determined by values $T(E_{1}), \dots, T(E_{n})$

$$\frac{P_{reg}}{P_{reg}} \quad \text{Let } X = X_{r}E_{1} + \dots + X_{n}E_{n}, \text{ be the unique expression of X
in terms of E_{1}, \dots, E_{n}

$$T(X) = T(X_{1}E_{1} + \dots + X_{n}E_{n}) = T(X_{1}E_{1}) + \dots + T(X_{n}E_{n}) = (\Gamma(X_{1}), \dots, T(E_{n}))$$

$$= X_{1} \cdot T(E_{1}) + \dots + X_{n}T(E_{n})$$

$$\frac{P_{reg}}{P_{reg}} = Let X = X_{r}E_{1} + \dots + X_{n}E_{n} = T(X_{n}E_{1}) = (\Gamma(X_{n}E_{n}) + \dots + X_{n}E_{n}) = T(X_{n}E_{1}) + \dots + T(X_{n}E_{n}) = (\Gamma(X_{n}) + \dots + X_{n}T(E_{n}))$$

$$\frac{P_{reg}}{P_{reg}} = Let X = V_{n}E_{1} + \dots + V_{n}E_{n} = T(X_{n}E_{1}) + \dots + T(X_{n}E_{n}) = (\Gamma(X_{n}) + \dots + X_{n}T(E_{n}))$$

$$\frac{P_{reg}}{P_{reg}} = Let X = V_{n}E_{1} + \dots + V_{n}E_{n} = T(X_{n}) + \dots + T(X_{n}) = (\Gamma(X_{n}) + \dots + X_{n}) = T(X_{n}) + \dots + V_{n}$$

$$\frac{P_{reg}}{P_{reg}} = Let X = V_{n} + V_{n} + V_{n} = T(X_{n}) + \dots + V_{n}$$

$$\frac{P_{reg}}{P_{reg}} = Let X = V_{n} + V_{n} + V_{n} = T(X_{n}) + \dots + V_{n} + V_{n}$$

$$\frac{P_{reg}}{P_{reg}} = Let X = V_{n} + V_{n} + V_{n} = T(X_{n}) + \dots + V_{n} + V_{n} = V_{n}$$

$$\frac{P_{reg}}{P_{reg}} = Let X = V_{n} + V_{n} + V_{n} = V_{n} + V_{n} + V_{n} + V_{n} = V_{n} + V_{n} + V_{n} + V_{n} + V_{n} = V_{n} + V_{n} +$$$$

Choose WI, W2,..., Wn E W.

Then there exists a unique lineas map T:V -> W s.t.

$$T(E_{1}) = w_{1}, T(E_{2}) = w_{2} \dots T(E_{n}) = w_{n}$$

$$P(co) = w_{1}coe^{2} \times = x_{1}E_{1} + \dots + x_{n}E_{n}$$

$$Degine T(x) = x_{1}w_{1} + \dots + x_{n}w_{n}$$

$$T_{1}s \text{ lineous } T(E_{1}) = w_{2}$$

• The module associated to a linear map

$$T: V \longrightarrow W \quad \text{encodes}$$

$$\{ \{ \{ \{ \}, \} \in G_{1}^{m} \} \mid j \} \in G_{1}^{m} \text{ Image} \}$$
basic for V basis for W .

$$T \text{ is completely descendence by (Expression f(eij) in earns of $\mathcal{I}_{1,m} \mathcal{I}_{m}$

$$T(E_{1}) = a_{11} \mathcal{I}_{1} + a_{21} \mathcal{I}_{2} + \dots + \mathcal{I}_{m} \mathcal{I}_{m}$$

$$T(E_{2}) = a_{12} \mathcal{I}_{1} + a_{22} \mathcal{I}_{2} + \dots + \mathcal{I}_{m} \mathcal{I}_{m}$$

$$I = T(E_{1}) = a_{1n} \mathcal{I}_{1} + a_{2n} \mathcal{I}_{2}$$

$$I.e: T(E_{2}) = \sum_{i=1}^{m} a_{ij} \mathcal{I}_{i} (j = 1, \dots, n)$$

$$f_{i} = \sum_{i=1}^{m} a_{ij} \mathcal{I}_{i} (j = 1, \dots, n)$$

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$$f_{i} = \sum_{i=1}^{m} a_{ij} \mathcal{I}_{i} (j =$$$$

.

$$T(E_{j}) = \sum_{i=1}^{n} a_{ij} \varphi_{i}$$

$$Q_{j}^{n} \qquad m(T) \sum_{i=1}^{n} = (a_{ij})_{1 \leq i \leq m}$$

$$I \qquad (i \leq j \leq m)$$

$$Matrix of T with respect to S on the left and I in the light.
$$m(T) \sum_{i=1}^{n} = matrix of T.$$

$$wre \leq \alpha_{i} eqt.$$

$$T: v \rightarrow w$$

$$I \qquad \alpha_{i} eqt.$$

$$J. \qquad basis \leq basis I$$$$

• Composition formula:
Suppose use have linear maps

$$\begin{array}{c} \mathcal{U}_{z} = \sum_{i=1}^{n} \mathcal{V}_{z} = \sum_{i=1}^{n} \mathcal{V}_{z$$

want to express Cite in terms e) (aij k) (bij).

$$(S \circ T) (E_{\kappa}) = S(T(E_{\kappa})) = S\left(\sum_{j=1}^{n} a_{j\kappa} e_{j}\right) =$$

$$= \sum_{j=1}^{n} a_{j\kappa} S(P_{j}) = \sum_{j=1}^{n} a_{j\kappa} \left\{\sum_{i=1}^{n} b_{ij} e_{ij}\right\} =$$

$$= \sum_{j=1}^{n} \sum_{i=1}^{n} a_{j\kappa} b_{ij} e_{ij} b_{ij} e_{ij} f_{i} =$$

$$(a_{j\kappa} b_{ij} = b_{ij} a_{j\kappa} e_{immation} e_{ij} e_{ij} e_{ij} f_{i} =$$

$$(a_{j\kappa} b_{ij} = b_{ij} a_{j\kappa} e_{immation} e_{ij} e_{ij} e_{ij} f_{i} =$$

$$(S \circ T) (E_{\kappa}) = \sum_{i=1}^{n} \left\{\sum_{j=1}^{n} b_{ij} a_{j\kappa}\right\} + \frac{\sum_{i=1}^{n} \left\{\sum_{j=1}^{n} b_{ij} a_{ij}\right\} + \frac{\sum_{i=1}^{n} \left\{\sum_{i=1}^{n} b_{ij} a_{ij}\right\} + \frac{\sum_{i=1}^{n} \left\{\sum_{i=1}^{n} b_{ij} a_{ij}\right\} + \frac{\sum_{i=1}^{n} \left\{\sum_{i=1}^{n} b_{ij} a_{ij}\right\} +$$

$$M(\tau) \stackrel{Q}{\underline{\pi}} = (b_{1j})_{1} \underbrace{i \cdot i \cdot \underline{\pi}}_{1 \cdot \underline{j} \cdot \underline{\beta}}_{1 \cdot \underline{\beta}}$$

Proposition Let
$$\Sigma = \{E_1, \dots, E_n\}$$
 basis for V .
 $\overline{\Psi} = \{\Psi_1, \dots, \Psi_n\}$.
 $M(Jd)_{\overline{\Psi}} = a_{ij}$ $M(Jd)_{\overline{\Sigma}} = \{b_{ji}\} = B$.
 $= A$

Then A, B are invertible and
$$B=A^{-1}$$

Proof: $M(Jd) \stackrel{\mathcal{Z}}{\leq} = M(Jd \circ Jd) \stackrel{\mathcal{Z}}{\leq}$
composition
 $= M(Jd) \stackrel{\mathcal{Z}}{\leq} M(Jd) \stackrel{\mathcal{Z}}{\leq}$

So
$$Id = AB$$

 $M(Ia)_{\underline{T}}^{\underline{T}} = M(Ia)_{\underline{S}}^{\underline{T}} M(Ia)_{\underline{T}}^{\underline{S}}$
 $In = BA$.
 $So AB = BA = I$.
 $V = Q^{3} = \int \begin{pmatrix} x_{1} \\ x_{2} \end{pmatrix} \quad x_{1} \in Q^{3}$.
 $\Xi = \int \begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix}, \begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix}, \begin{pmatrix} 0 \\ 0 \\ 1 \end{pmatrix} \end{pmatrix} \qquad basis$
 $\overline{I} = - \int \begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix}, \begin{pmatrix} 0 \\ 1 \\ 0 \end{pmatrix}, \begin{pmatrix} 0 \\ 1 \\ 1 \end{pmatrix} \qquad basis$
Take $T: Q^{3} \longrightarrow Q^{3}$ to be. Commute $M(T)_{\underline{S}}^{\underline{S}}$
 $T(x_{1}) = \begin{pmatrix} 2x_{1} + x_{2} - x_{3} \\ 3x_{2} + x_{3} \\ -3x_{2} + x_{3} \end{pmatrix}$
 T is lineas (he knows it).

$$T\begin{pmatrix} i \\ 0 \\ 0 \end{pmatrix} = \begin{pmatrix} i \\ 0 \\ 0 \end{pmatrix}$$

$$T\begin{pmatrix} 0 \\ i \\ 0 \end{pmatrix} = \begin{pmatrix} 1 \\ 3 \\ 0 \end{pmatrix}$$

$$M(T) \underset{\mathcal{S}}{\mathcal{S}} = \begin{pmatrix} z & i & -i \\ 0 & 3 & i \\ 0 & 0 & 4 \end{pmatrix}$$

$$F_{in} d = M(T) \underset{\mathcal{G}}{\underline{\mathcal{G}}}$$

$$T\begin{pmatrix} 0 \\ i \\ 0 \end{pmatrix} = \begin{pmatrix} -i \\ 4 \\ 4 \end{pmatrix}$$
Use composition between

$$T = Id \circ T \circ Id .$$

$$M(T)_{\Psi}^{\Psi} = M(Id)_{\Sigma}^{\Psi} M(T \circ Id)_{\Psi}^{\Sigma} =$$

$$= M(Id)_{\Sigma}^{\Psi} m(T)_{\Sigma}^{\Sigma} M(Id)_{\Psi}^{\Sigma}$$

$$kncw \quad M(T)_{S}^{\Sigma}$$
also know
$$M(Id)_{\overline{\Sigma}}^{\Sigma} = \begin{pmatrix} 1 & 1 & 0 \\ 0 & 1 & 1 \\ 0 & 0 & 1 \end{pmatrix}$$

$$Check inverse \quad og \quad = \begin{pmatrix} 1 & -1 & 1 \\ 0 & 1 & -1 \\ 0 & 0 & 1 \end{pmatrix} = M(Id)_{\overline{\Sigma}}^{\overline{\Sigma}}$$

50 now.

$$\begin{split} \mathcal{M}(T) \overset{\Phi}{\mathbf{p}} &= \begin{pmatrix} 1 & -1 & 1 \\ 0 & 1 & -1 \\ 0 & 0 & 1 \end{pmatrix} \cdot \begin{pmatrix} 2 & 1 & -1 \\ 0 & 3 & 1 \\ 0 & 0 & 4 \end{pmatrix} \begin{pmatrix} 1 & 1 & 0 \\ 0 & 1 & 1 \\ 0 & 0 & 1 \end{pmatrix} = \\ &= \begin{pmatrix} 2 & -2 & -2 \\ 0 & 3 & -3 \\ 0 & 0 & 4 \end{pmatrix} \begin{pmatrix} 1 & 1 & 0 \\ 0 & 1 & 1 \\ 0 & 0 & 1 \end{pmatrix} = \begin{pmatrix} 2 & 0 & 0 \\ 0 & 3 & 0 \\ 0 & 0 & 4 \end{pmatrix} \end{split}$$

November 29th

$$T: V \longrightarrow W$$
 enear
Basis $\Sigma = [E_1, ..., E_n] Jor V$
 $\overline{\Psi} = [\Psi_1, ..., \Psi_m] for W$

$$T(E_{j}) = \sum_{i=1}^{\infty} a_{ij} P_{i}$$

$$Define \quad M(T) = (a_{ij}) \quad i \leq i \leq m \quad Matriz = \left(\begin{array}{c} P_{i} & P_{i} & P_{i} \\ P_{2} & P_{2} & P_{2} \end{array} \right)$$

$$I \leq j \leq n \quad P_{2} \quad P_{2}$$

Linear maps
$$\longrightarrow$$
 matrices
 $T \longrightarrow m(T) \stackrel{E}{=}$
Conversely, y (aij) is is m
 $i \le j \le m$
determine linear map: $T \cdot V \rightarrow W$
by $T(E_j) = \sum_{i=1}^{m} a_{ij} \varphi_i$

This gives a origine linear map. In this case: $M(T) \stackrel{\text{$!}{$!}}{\underset{\text{$!}{$!}}{$!}} = A$.

Under this correspondence

$$\begin{cases} Composizion of lineas \\ maps \\ \end{cases} \begin{cases} <-> \\ product \\ product \\ product \\ \end{cases} \\ (M (S \circ T)_{\mathcal{E}}^{\mathbf{F}} = M(S)_{\mathbf{F}}^{\mathbf{F}} = M(T)_{\mathcal{E}}^{\mathbf{F}} \end{cases}$$

As composition is Associative then matrix product is also associative. Standard example ((Again) (et $A = (aij)_{j \in i \in M}$ $j \in j \in N$.

In this case
$$M(T) = A$$

 $T_{n} : IF^{n} \rightarrow F^{m}$.
 $T_{n}(S) = A_{\Sigma}$ matrix product
Prop $M(T_{n}) = A$
 $Chack:$
 $T_{n}(F_{n}) = \begin{pmatrix} a_{n} & a_{n} & \dots & a_{n} \\ a_{21} & a_{22} & \dots & a_{2n} \\ \vdots & \vdots & a_{m} & a_{m_{Z}} & \dots & a_{m_{N}} \end{pmatrix} \begin{pmatrix} i \\ 0 \\ i \\ 0 \end{pmatrix} = \begin{pmatrix} a_{11} \\ a_{21} \\ a_$

$$\begin{aligned} \text{Coloniant} \quad \mathbf{M}(\mathbf{G})_{\mathbf{Z}}^{\mathbf{Z}} \\ O \left(exp(x)\right) &= exp(x)^{\frac{1}{2}} \varphi(xexp(x) + \Theta \cdot x^{\frac{1}{2}} exp(x)) \\ D \left(x exp(x)\right) &= exp(x)^{\frac{1}{2}} \varphi(xexp(x)) \\ D \left(x^{\frac{1}{2}} exp(x)\right) &= 2x \cdot exp(x)^{\frac{1}{2}} xexp(x) \\ \mathbf{M}(\mathbf{D})_{\mathbf{Z}}^{\mathbf{Z}} &= \mathbf{M}(\mathbf{D})_{\mathbf{Z}}^{\mathbf{D}} \quad \mathbf{M}(\mathbf{D})_{\mathbf{Z}}^{\mathbf{D}} &= \begin{pmatrix} 1 & 1 & 0 \\ 0 & 1 & 2 \\ 0 & 0 & 1 \end{pmatrix} \\ \text{Suppose I take } \mathbf{D}^{\frac{1}{2}} = \mathbf{D} \circ \mathbf{D} \circ \mathbf{D} \cdot \begin{bmatrix} 1 & 1 & 0 \\ 0 & 1 & 2 \\ 0 & 0 & 1 \end{bmatrix} \\ \text{New take } \mathbf{D}^{\frac{1}{2}} = \mathbf{M}(\mathbf{D})_{\mathbf{Z}}^{\mathbf{D}} \quad \mathbf{M}(\mathbf{D})_{\mathbf{Z}}^{\mathbf{E}} &= \begin{pmatrix} 1 & 1 & 0 \\ 0 & 1 & 2 \\ 0 & 0 & 1 \end{pmatrix} \\ \text{New take } \mathbf{D}^{\frac{1}{2}} = \mathbf{M}(\mathbf{D})_{\mathbf{Z}}^{\mathbf{E}} \quad \mathbf{M}(\mathbf{D})_{\mathbf{Z}}^{\mathbf{E}} &= \begin{pmatrix} 1 & 1 & 0 \\ 0 & 1 & 2 \\ 0 & 0 & 1 \end{pmatrix} \\ \text{New take } \mathbf{D}^{\frac{1}{2}} = \mathbf{M}(\mathbf{D})_{\mathbf{Z}}^{\mathbf{E}} \quad \mathbf{M}(\mathbf{D})_{\mathbf{Z}}^{\mathbf{E}} &= \begin{pmatrix} 1 & 1 & 0 \\ 0 & 1 & 2 \\ 0 & 0 & 1 \end{pmatrix} \\ \text{New take } \frac{d^{\frac{1}{2}}}{dx^{2}} - \frac{d^{\frac{1}{2}}}{dx^{2}} + \frac{d}{dx} \end{pmatrix} \mathbf{Q}(x) = x \cdot exp(x) - 2 \cdot x^{\frac{1}{2}} \cdot exp(x) \\ \text{M}(\mathbf{D}^{\frac{1}{2}} - \frac{d^{\frac{1}{2}}}{dx^{2}} - \frac{d^{\frac{1}{2}}}{dx^{2}} + \frac{d}{dx} \end{pmatrix} \mathbf{Q}(x) = x \cdot exp(x) - 2 \cdot x^{\frac{1}{2}} \cdot exp(x) \\ \text{M}(\mathbf{D}^{\frac{1}{2}} - \mathbf{D}^{\frac{1}{2}} + \mathbf{D}^{\frac{1}{2}} - \mathbf{D}^{\frac{1}{2}} \\ \frac{1}{0} \circ \mathbf{D}^{\frac{1}{2}} + \frac{1}{0} \\ \frac{1}{0}$$

$$(DII) \sim \begin{pmatrix} 1 & 1 & 0 \\ 0 & 1 & 2 \\ 0 & 0 & 1 \end{pmatrix} \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix}$$
$$D = \begin{pmatrix} 1 & -1 & 2 \\ 0 & 0 & -2 \\ 0 & 0 & 1 \end{pmatrix}$$
$$D = \begin{pmatrix} 1 & -1 & 2 \\ 0 & 0 & -2 \\ 0 & 0 & 1 \end{pmatrix} \begin{pmatrix} 2 & 0 & 7 \\ -7 & 0 \\ 0 & 0 & 1 \end{pmatrix}$$
$$\int \begin{pmatrix} 2 & 0 & 7 \\ -7 & 0 \\ 1000 \end{pmatrix} = \begin{pmatrix} 1 & -1 & 2 \\ 0 & 0 & 1 \end{pmatrix} \begin{pmatrix} 2 & -7 \\ 1000 \end{pmatrix} = \begin{pmatrix} 2 & 0 & 7 \\ -7 & 0 \\ 1000 \end{pmatrix}$$
$$= 2007 \exp(k) - 2005 \times \exp(k) + 1000x^{2} \exp(k) .$$
$$T.V \longrightarrow W \quad Intesc.$$
$$Degine \quad Her(T) = \begin{cases} x \in V : T(x) = 0 \\ y \\ T \end{pmatrix}$$
$$ker(t) .$$
$$Ker(t) = \begin{cases} x \in V : T(x) = 0 \\ y \\ T \end{pmatrix}$$
$$ker(t) .$$
$$Reposition \quad ker(T) \subset V \quad ker(T) \text{ is a vector spice}.$$
$$Red \quad v = ker(T) \quad y \in ker(T) = x + y \in ker(T) \\ T(x+y) = T(x) + T(y) = 0 + 0 = 0 .$$
$$T(x+y) = T(x) + T(y) = 0 + 0 = 0 .$$

$$T(\lambda x) = \lambda t(x) = \lambda \cdot 0 = 0$$

=> $\lambda \cdot x \in \text{ker}(t)$.

The remaining arisins are automatically satisfied because they are true already in V. November 30th 2018

Let V be a vector space / IF. and ket
$$U \in V$$

Sug U is a vector subspace of V when:
A) $0 \in U$.
A) $T_{3} \neq J_{1} \notin U$ then $\pm J_{2} \in U$.
A) $T_{3} \neq J_{1} \notin U$ then $\pm J_{2} \in U$.
A) $T_{3} \neq U \in V$ is a vector subspace then:
At is they a vector space.
Proof we have all necessary elements of proverve and
ell avients satisfies because already subspace in V.
Refinition: Let T: V -> W be there, define:
Keinel of T.
Required: Ker(T) = { $x \in V + T(x) = 0$ }
Required: Ker(T) is a electric subspace of V(= domain of T).
Proof : E) $0 \in \text{ker(T)}$.
T(x) = 0 T(y) = 0
So $T(x + j) = T(x) + T(y) = 0 + 0 = 0$.
A) $f_{3} \neq e = \text{ker(T)}, \lambda \in F$.
 $T(x_{1}) = NT(x) = \lambda \cdot 0 = C$.
 $\lambda x \in \text{ker(T)} \in O$.
Requires $T(x) = N = 0$.
A) $T_{3} \neq e = \text{ker(T)} \oplus O = 0$.
A) $T_{3} \neq e = \text{ker(T)} \oplus O = 0$.
A) $T_{3} \neq e = \text{ker(T)} \oplus O = 0$.
 $T(x_{1}) = NT(x) = \lambda \cdot 0 = C$.
 $\lambda x \in \text{ker(T)} \oplus O = 0$.
 $Cogne = Tn(T) = { w \in W, \exists V \in V = T(v) = w Y = 0$.
 $T(x_{2}) = T(x) = T(x) = 0$.
 $T(x_{2}) = T(x) = 0$.
 $T(x_{3}) = NT(x) = x \cdot 0 = C$.
 $T(x_{3}) = NT(x) = x \cdot 0 = C$.
 $T(x_{3}) = NT(x) = x \cdot 0 = C$.
 $T(x_{3}) = NT(x) = 1 = 0$.
 $Cogne = Tn(T) = { w \in W, \exists V \in V = T(v) = w Y = 0$.
 $T(x_{3}) = T(x_{3}) = T(x_{3}) = T(x_{3}) = 0$.
 $T(x_{3}) = T(x_{3}) = T(x_{3}) = T(x_{3}) = 0$.

Proposition If
$$T: V \longrightarrow W$$
 (near than $Im(T)$ is a vector
Subspace of W .
Freed A) $O_W \in Im(T)$ because $T(O_V) = O_W$.
Fix $(V_1, W_2) \in Im(T)$. Choose $U_1, V_2 \in V$.
S.F. $T(V_1) \supset W$, $T(V_2) \supset W_2$
 $T(V_1 + V_2) = T(V_1) + T(V_2) = W_1 + W_2$
F($V_1 + V_2$) = $T(V_1) + T(V_2) = W_1 + W_2$
F($V_1 + V_2$) = $T(V_1) + T(V_2) = W_1 + W_2$
F(V_1) = $A T(V) = A W Se A W \in Im(T)$
Theorem: Kennel - Ranke Theorem
 $T(AU) = A T(V) = A W Se A W \in Im(T)$
Theorem: Kennel - Ranke Theorem W .
 $We that A dim kee(T) + dim Im(T) = dum V$.
 $We that A dim kee(T) + dim Im(T) = dum V$.
 $We that A source for the form Im(T) = dum V$.
 $We that A source for the form Im(T) = dum V$.
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 $We there form Im(T)$.
 $We there for the form Im(V)$.
 $We there form Im(T)$.

Chain
$$\{E_1, \dots, E_{k+m}\}$$
 is a share for V .
 $\{E_i, \dots, E_{k+m}\}$ is $L \equiv .$
Suppose $\{h, E_i + \lambda_i, E_i + \dots + \lambda_k E_{k+1} + \lambda_{k+m} E_{k+m}\} = 0$
 $\{i \in F\}$
 I have to show each $\lambda_i = 0$.
 $h(r) = T(\lambda_i (E_i + \dots + \lambda_k E_k + \lambda_{k+1} E_{k+1} + \dots + \lambda_{k+m} E_{k+m}) = 0$
 $h(r) = T(E_i) + n = \lambda_k E_k + \lambda_{k+1} E_{k+1} + (E_{k+1}) + \dots + \lambda_{k+m} T(E_{k+m}) = 0$
 $h(r) = K_{k+1} + (E_{k+1}) + \lambda_{k+1} + (E_{k+1}) + \dots + \lambda_{k+m} T(E_{k+m}) = 0$
 $i \in \lambda_{k+1} + (E_{k+1}) + 1 + h_{k+m} T(E_{k+m}) = 0$
 $i \in \lambda_{k+1} = (E_{k+1}) + 1 + h_{k+m} T(E_{k+m}) = 0$
 $i \in \lambda_{k+1} = (E_{k+1}) + 1 + h_{k+m} T(E_{k+m}) = 0$
 $i \in \lambda_{k+1} = h_{k+2} = \dots = h_{k+m} = 0$
 $h_{k+1} = h_{k+2} = \dots = h_{k+m} = 0$
 $h_{k+1} = h_{k+2} = \dots = h_{k+m} = 0$
 $h_{k+1} = h_{k+2} = \dots = h_{k+m} = 0$ find
 $h_{k+1} = h_{k+2} = \dots = h_{k+m} = 0$ find
 $h_{k+1} = h_{k+2} = \dots = h_{k+m} = 0$ for $a \leq i$
 $i \in i \leq k+m$
 $i \in \{E_{i}, \dots, E_{k+m}\} = h + I$
 $Wr also comm that $\{E_{i}, \dots, E_{k+m}\}$ spans V .
 $(E_{k+1} \leq k + M \in I_{k+1})$
 $h_{k+1} = h_{k+2} = I_{m} = 1$
 $k \in V$ $(I) we got to show $\lambda = x_i E_i + \dots + I_{k+m} E_{k+m}$ for
 $Some x_{1}, \dots, X_{k+m} \in I_{k}$
 $\cdot Considen T(X) = I_{m}(\tau)$.
 $and express T(X) in terms of P apply Γ .
 $\{P_{1}, \dots, P_{m}\}$$$$

$$T(x) = \xi, \quad (f_1 + \dots + f_m \quad f_m \quad f_m \quad d_m \quad$$

Here
$$T(x) = 0$$
 and $T(0) = 0$.
Since T is injective $T(x) = 0$ = $0 \Rightarrow ker(T) = 10^{1}$.
Since T is injective $T(x) = 0 \Rightarrow ker(T) = 10^{1}$.
So oppose $ker(T) = 10^{1}$ and that $T(k) = T(k^{1})$.
SP $T(x-x^{1}) = T(k) - T(k^{2}) = 0 \Rightarrow x-k^{2} = 0$ B
 $x = x^{1}$ and T is injective GED.
Symmetric $ker(T) = \frac{10^{1}}{10^{1}}$
Let $E_{1,r_{1}}E_{n}$ be a basis for V .
Goins $T(E_{1}), \dots, T(E_{n})$ is a solid for $J(T)$.
Let Ψ e $T(n(T)$.
Choose $x \in V + T(\Sigma) = W$.
Where $x = x, E_{1} + \dots + x_{n} E_{n}$.
 $\Psi = T(x) = x, T(E_{1}) + \dots + x_{n} T(E_{n})$
so $T(E_{1}), \dots, T(E_{n})$ spaces $T(n(T)$.
For $\lambda_{1}T(E_{1}) = \dots + \lambda_{n} E_{n} = 0$.
 $T(\lambda_{1}E_{1}) = \dots + \lambda_{n} E_{n} = 0$.
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 $f(\lambda_{1}E_{1} + \dots + \lambda$

 $\underbrace{example}_{A=} A_{=} \begin{pmatrix} 1 & 1 & -1 & 1 & 1 & -1 \\ 1 & 1 & 1 & -1 & 1 & -1 \\ 1 & 1 & -3 & 3 & 1 & -1 \end{pmatrix}$ over Q $T_A: Q^6 \longrightarrow Q^3$ $T_A(x) = A \cdot X$ So $ker(T_A) = solution set of homogeneous system AX.$ Reduce A to echelon $\begin{pmatrix} 1 & 1 & -1 & 1 & -1 \\ 0 & 0 & 2 & -2 & 0 \\ 0 & 0 & -2 & 2 & 0 \end{pmatrix} \longrightarrow \begin{pmatrix} 1 & 1 & -1 & 1 & 1 & -1 \\ 0 & 0 & 1 & -1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{pmatrix}$ $\begin{pmatrix} x_{1} \\ x_{2} \\ x_{3} \\ x_{4} \\ x_{5} \\ x_{6} \\ x_{7} \\ x_$ con la sol del sistema cuando prof X am O General solution $\begin{pmatrix} -x_2 & -x_5 + x_6 \\ x_2 \\ x_3 \\ x_4 \\ x_5 \\ x_6 \end{pmatrix} = \begin{pmatrix} -x_5 + x_6 \\ -x_6 \\ -$ COMMON MISTARES Take the columns above curiled variables in REDUCED Matrix In this case $\begin{pmatrix} 1 \\ 0 \\ 0 \end{pmatrix} \begin{pmatrix} 0 \\ 1 \\ 0 \end{pmatrix}$ CORRECT METHOD Take columns above wicked variables m ORIGINAL MATRIX $\begin{pmatrix} 1 \\ 1 \\ 1 \\ -3 \end{pmatrix}$; bossis for Im In this case

(41)

Examplez :

$$A = \begin{pmatrix} 2 & | & | & | & | \\ 3 & 2 & - | & | & | \\ 1 & | & -2 & 0 & 0 \end{pmatrix}$$

and reduce. $\begin{pmatrix} 1 & | & -2 & | & 0 & | \\ 0 & -1 & 5 & | & | \\ 0 & -1 & 5 & | & | \end{pmatrix} \longrightarrow \begin{pmatrix} 1 & | & | & | & | \\ 0 & | & 0 & | & 0 & 0 \end{pmatrix}$
$$S_{0} \begin{pmatrix} 2 \\ 3 \end{pmatrix} \begin{pmatrix} 1 \\ 2 \end{pmatrix} : it a basis for Im(T_{n})$$

$$S_{0}probe
$$\begin{cases} a_{11} \times 1 + \dots + a_{nn} \times n = b_{1} \\ a_{n1} \times 1 + \dots + a_{nn} \times n = b_{n} \\ a_{n1} \times 1 + \dots + a_{nn} \times n = b_{n} \\ a_{n1} \times 1 + \dots + a_{nn} \times n = b_{n} \\ A_{n1} \times 1 + \dots + a_{nn} \times n = b_{n} \\ A_{n1} \times 1 + \dots + a_{nn} \times n = b_{n} \\ A_{n1} \times 1 + \dots + a_{nn} \times n = b_{n} \\ A_{n2} \times 1 + \dots + a_{nn} \times n = b_{n} \\ S_{0} \text{ suppose } A_{2} = b \\ S_{0} \text{ suppose } A_{2} = b \\ S_{0} \text{ suppose } A_{2} = b \\ A_{2} - A_{2} - b \\ M_{2} - A_{2} - b \\ M_{2} - M \\ M_{2$$$$

· BASIS THEOREM

let V be a non-zero vector space over field IF. 1) There exists at least one basis for V. (EXISTENCE) 2) Any two basis for V have same no og elementes (= dm (V)) (UNIQUENESS) Proposition: Let (W, ..., why C V V is a vector space / IF. If O E TWILL WAY then Yw, ..., wa) is NOT LE. Proof Suppose wit = 0 (uno de esos vectores es 0) consider $\sum_{j=1}^{\infty} \lambda_j w_j$, where $\lambda_j = \int_{-\infty}^{\infty} j = i - \sqrt{\lambda_i} w_i = \lambda_i c \Rightarrow \lambda_i \neq 0$ then $\sum_{j=1}^{2} \lambda_j w_j = 0$ but $\lambda_i \neq 0$ so $\gamma w_1, \dots, w_n \gamma$ is not LI $d_i \cdot w_i + \lambda_j w_j = 0$ Proposition: Let V be a nonzero versor opace spanned by vertors Yw, ..., way Then V has a basis with at nost a elements. Proof By inducción on n Let P(n) be statement of the proposition N=1 · First prove P(1). → V≠0 is spanned by a single element w = (w,) () I claim that Ywy is L. J. · Suppose NW =0 Got to show $\lambda = 0$. - Suppose not. Then multiply across by h-1. x-' x 20 00 - 20 - 0 = 0. + But quy & spans V. so V=0 Contraduction, so P(1) is true.

+ Suppose
$$F(n_{i})$$
 is true.
Got to show $P(n)$ is true.
V is spanned by $\int U_{i_{1}}, \dots, U_{i_{n}} Y$.
Egg $\{U_{i_{1}}, U_{i_{n}}, U_{i_{n}} Y = 0$, $U_{i_{n}} = 0$, $U_{$

If you believe Axion of CHOLCE, then every non-zero vector year.
has a basis regardless if whether it is finitely generated of not.
If not then Fills.
IMPROVENESS
Recay by E. Stein the c 1908
Exclarge connec (Basy verien) to Suppose
$$\{w_1, \dots, w_n\}$$
 open V.
and let $y \in V$ [$y \neq 0$]
Write $y = \sum_{i=1}^{n} \lambda_i w_i$
 $i \neq i$
 $(i \in . Swap ψ_i and ψ_i
 $i \neq i$
 $(i \in . Swap ψ_i and ψ_i
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 $(i = . Swap ψ_i and ψ_i
 $i = (i + .) \psi_i + ... (j \neq i)$
 $(i = . Swap ψ_i
 $i \neq i$
 $i \neq i$$$$$$$$$$$$$$$$$$$$

Full exchange lemma

Let V be a non zero vector space spanned by {
$$w_1, \ldots, w_n$$
 y
Suppose Y_{u_1}, \ldots, y_k y is L. I ($y_i \in V$).
Then:
i) k $\leq n$.
ii) There exists a spanning set $(w_{i_1}^{i_1}, \ldots, w_n$ y for $R \in I \leq j$.
s.t $w_i' = v_i$ ($1 \leq i \leq k$) $w_j' \in \{w_{i_1}, \ldots, w_n$ y for $R \in I \leq j$.

December 7th 2018

Exchange Lemma (Stents, 1908):
Let
$$\{i \forall i, \dots, \forall n \forall i \in paan vector islane V.$$

and let $\{i \forall i, \dots, \forall k'\} \in V$ be LIT. Then:
L') $k \leq n$.

spans V.

$$I_{j} = \sum_{i=1}^{N} \mathcal{M}_{j} \mathcal{W}_{j} = \mathcal{M}_{i} \mathcal{W}_{i} + \sum_{j \neq i}^{j} \mathcal{M}_{j} \mathcal{W}_{j}$$

$$= \left(\frac{\mathcal{M}_{i}}{\lambda_{i}}\right) \mathcal{V}_{i} + \sum_{j \neq i}^{j} \left(\mathcal{M}_{j} - \mathcal{M}_{i} - \frac{\lambda_{j}}{\lambda_{i}}\right) \mathcal{W}_{j}$$

Te we've proved induction base.

· Induction step: Assume true for k-1

where
$$k-1 \leq n$$
. $w_j \in j w_1, ..., w_n \gamma$.
In this case $k-1 \leq n$.

otherwise if n=k-1 then yu, ..., Vk-1] spans V.

· This gives a contraduction by expressing

why?
why?
where
$$y = k-1$$
 then $yy_{1,...,y_{k-1}}$ spans V .
is gives a contradiction by expressing
 $V_{k} = \sum_{i=1}^{k-1} \eta_{i} V_{i} \eta_{k} + \int_{V_{i}} \int_{V_{$

÷.

· As k-1 < n then k < n, which is pirst part of conclusion. 6 ¹ To given, take spanning set gui, Vk-1; Wk, ..., Why

· Express Vic in terms of Write $\hat{\langle}$

$$V_{k} = \lambda_{i} u_{i} t + \lambda_{k-1} v_{k-1} + \sum_{j=k} \lambda_{i} w_{i}$$

 $Pb_{Rive} V_{k} \neq 0$ as $\gamma v_{i} \cdots v_{k} \gamma v_{k} \gamma$ is L.J.

· Claim that by \$ 9 for some j. KEJEN ... Z · otherwise, if hi = o be KEILA I have a dependence relation. VR = hivit ... + hzvz proque deboreater Q = hvit ... + hzvz Contradicting lineas independence of VI, VKY. Vir=huit + ... + hk-1 Vir-1 + hi wit + 5 hi wi hi to j=k hi wit + 5 hi wi hi to j=k hi wit + Shi wit (countries with hi to j=k hi wit + Shi with to him him he (countries with the him him he has he h So hito for some j. KLiLn · So use baby version and swap vie and wi New we have a spanning set of form YVi, where, VKY U Juj; K Li En Y cardinaly (n-k) hb of elements. Formally to complete proof: Put $w_i'' = V_i$ I Likk and let $\int w_j'' : k+1 \leq j \leq n \int$ be some indering of the set dw_j . 15 j 5 k). Q.ED Corollasy BASIS THEOREM . Let V be a non-zero vector space, then: >>> V has at least one basis ii) Any two basis have same no of elements. Proof: i) already done. ic) Suppose YE, ..., Emy is a basis for V. and Y 9, ..., Phy is also a basis /V of E, ..., Eny is L. I and of Q, ..., Q'ny spons . So min. Now de qui a y is LI (EI, En y spans, so nem.

SO MENEM SO MEN. QED

(crollary (of Exchange lemme) -Suppose: V, W are rector spaces with VCW. then, $V = W \iff \dim(V) = \dim(W)$. Proof => Trivial <= Suppose, dim (v) = dim (w) 1 · IJ dim (v) = dim (w)=0 then V=104=w. So rothing to prove. · So suppose V to (so w to). Let YE, Eny be a basis for V. let der, Pry beabasis for w. 19, and 9, 9 spans W. JE, ..., EngisLII By exchange Lemme: YE, En Y spans W. IJ W E W . wite . $U = \lambda_{i} \in (+, +, +, \lambda_{n} \neq N)$ because each Eie V. So wevew So V=W QED

Back to kenel - Rank Theorem:

$$T: V \rightarrow W$$
 be ineas
 $dim(v) = dim ker(T) + dim Im(T)$

I slowed T injective (=> dim ker(t) =0.

Proposition: T is surjective (=> dim Jm(T) = dim (w).

Proos .

Trivial

$$F = A_S Jm(T) = \omega$$
 then dim $Jm(T) = \dim(\omega)$
 $Jmplies Jm(T) = \omega$ J.e. T surjective

· Isomorphisms . Let V, W be vector spaces Say that V, W are isomorphic written VZW when there exists a bjective linear map t: V --> W. 1 1) V~ V (Reflexivity) Idy V -> V lineas and sigerive 2) IJ V = W then W = V (Symmetry) Proof let T: V -> W be linear and bijective I claim T-1: W-> V is linear (and it is certainly bijective) Let W, WZ EW. Consider $T^{-1}(\omega_{1}+\omega_{2}) - T^{-1}(\omega_{1}) - \Gamma^{-1}(\omega_{2})$ $\mathcal{T}(\tau^{-1}(\omega_1 + \omega_2) - \tau^{-1}(\omega_1) - \tau^{-1}(\omega_2)) = \tau^{-1}(\omega_1 + \omega_2) - \tau^{-1}(\omega_1) - \tau^{-1}(\omega_2) =$ Apply T $= w_1 + w_2 - w_1 - w_2 = 0$ But + is injective, so: $T^{-1}(\omega_1 + \omega_2) - T^{-1}(\omega) - T^{-1}(\omega_2) = 0$ i.e. $T^{-1}(w_1 + w_2) = T^{-1}(w_1) + T^{-1}(w_2)$ Likewise $T(T^{-1}(\lambda\omega) - \lambda T^{-1}(\omega)) = 0$ T is injective, so $t^{-1}(\lambda w) - \lambda t^{-1}(w) = c r \delta e_r$ $T^{-1}(\lambda w) = \lambda T^{-1}(w)$. Ord T^{-1} is linear. QEO// (Symmetry). -3) MEV and VEW, then MEW. Transitivity Proop Ig T. M->V is linear bijective, S. V -> W is linear and by pereve Then sot : 11 -> 11 is linear bijecture

$$Ty \quad V, \ w \ are \ vector \ spaces / F, then:
V = W \quad C = 2 \dim(V) = \dim(W) .
$$Soppose \quad T: V = 0 \quad (webs \ and \ b)ectore
(et \quad 9h, ..., fn T \quad be a basis for V.
Charin relact $\{T(f_{1}), ..., T(f_{n})\}$ is a basis for W.

$$\{T(f_{n}), ..., T(f_{n})\} \quad is \quad L.T. .
Why??
$$\sum_{i=1}^{n} \lambda_{i} T(f_{n}) = 0$$

$$\lim_{i \neq 1} T\left(\sum_{i=1}^{n} \lambda_{i} F_{n}\right) = 0$$
But T is injective to kee(t) = 0; so

$$\lim_{i \neq 1} \lambda_{i} F_{n} = 0 \quad and \quad \{F_{1}, ..., F_{n}Y \ is LT, so \quad \lambda_{1} = ... = h = 0.$$
(Laim $T(F_{n}), ..., T(f_{n})$ spans W.
(Laim $T(F_{n}), ..., T(f_{n})$ spans W.
(Laim $T(F_{n}), ..., T(f_{n})$ spans W.
(Laim $T(F_{n}) = 1 \left(\sum_{i \neq 1}^{n} \lambda_{i} F_{i}\right)$

$$\lim_{i \neq 1} \frac{2}{m} \lambda_{i} F_{i} = 1$$

$$\lim_{i \neq 1} \frac{2}{m} \lambda_{i} F_{i}$$

$$\lim_{i \neq 1} \frac{2}{m} \sum_{i \neq 1} \lambda_{i} F_{i}$$

$$\lim_{i \neq 1} \frac{2}{m} \sum_{i \neq 1} \lambda_{i} F_{i}$$

$$\lim_{i \neq 1} \frac{2}{m} \sum_{i \neq 1} \lambda_{i} T(F_{i})$$

$$\lim_{i \neq 1} \frac{2}{m} \sum_{i \neq 1} \lambda_{i} T(F_{i})$$

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$$\lim_{i \neq 1} \frac{2}{m} \sum_{i \neq 1} \lambda_{i} T(F_{i})$$

$$\lim_{i \neq 1} \frac{2}{m} \sum_{i \neq 1} \sum_{$$$$$$$$

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.

Conversely,
$$g \dim (v) = \dim(w)$$

then $v \equiv w$.
Take basis $f \in 1, ..., Enf for v$.
 $(q_{1,...,}, q_{n} f for w)$.
Let $T : v \to w$ be the unique lines mapping determined by
 $T(E_{i}) = q_{i}$.
Let $s : w \to v$ be unique lines $s : t : S(q_{i}) = E_{i}$.
Then, $S_{0}T(E_{i}) = E_{i}$ so $S_{0}T = TdV$.
Likewise, $(T_{0}S)(q_{i}) = q_{i}$ so $T_{0}S = Edw$.
so $T : V \to w$ is a uncer byzerion
 $T^{-1} = S$.
 $w V \cong w$.
 $QED//$

December 15:
PERMUTATIONS

$$T_{n} = \langle \sigma : \{1, ..., n \rangle \} \longrightarrow \{1, ..., n \rangle$$

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 $T_{n} = \langle \sigma : \{1, ..., n \rangle \} \longrightarrow \{1, ..., n$

$$F_{3} = c = C_{n} \quad define \quad d_{n}(\sigma) = \overline{v} = \sigma(s) - \sigma(s)^{n}$$

$$= c + s = \overline{v}$$
so $d_{n}(r+1) = \overline{v} = (s - \tau)^{n}$

$$= c + s = \overline{v}$$
Production $T_{1} = (i, 1+i)$ is an adjust transplaced interplaced interval $f_{n}(\sigma t) = -d_{n}(\sigma)^{n}$
Prove $(i, 1+i)$ are gived
$$(r, s) = -d_{n}(\sigma)^{n}$$
Prove $(i, 1+i)$ are gived
$$(r, s) = \sigma(s) + \sigma(s)^{n}$$
Prove $(i, 1+i) = i^{n}$

$$(o) = (r, s) = \sigma(s)^{n} + i^{n} + i^{n}$$
Prove $(o) = (i, 1+i) = i^{n} + i^{n}$

$$(c) = (i, 1+i) = (i+i)^{n}$$

$$(c) = (i, 2i)^{n}$$

$$(c) = (i, 2i)^{n}$$

$$(c) = (i, 2i)^{n}$$

$$(c) = (c)$$

Fig. By induction on A.
A=1 Take
$$\sigma = z \circ , f_{i}(z) = -f_{i}(zd)$$
 by above.
So suppose process for A⁻¹
 $f_{i}(z_{1},...,z_{A_{i}},z_{A}) = -f_{i}(T_{i},...,T_{A_{i}}) = (-1) \cdot (-1)^{A-1} = (-1)^{A}$
By natures
 $f_{i}(z) = (-1)^{A}$ if T is a
 $f_{i}(z)$ product of A adjacent transported
II
 $g_{i}(z)$
December 14th
Dynamic of required new extract point
 $A = (a_{ij})_{i \leq j \leq A}$. it now
 $A = (a_{ij})_{i \leq j \leq A}$. it now
Condition 1
IJ $A_{i,k} \neq 0$ and $A_{i,k} \geq 0$, then $i \leq K$.
In English, zero ences care cost.
Let $A_{i,k}, ..., A_{i,k}$
be there zero tows.
For $i \leq i \leq T$.
 $c(i) = Min d j i a_{i,j} \neq 0$ 1 .
 $i \in ..., c(i)$ is the i^{T} course is unknew you get a not zero entry in
 $A_{i,k}$.
 $a_{i,j}(z)$ for but $a_{i,j} = 0$ if $j = c(i)$.
Condition 2 $a_{i,j}(z) = 1$.
Condition 3 $a_{k,j}(k) = 0$ if $k \neq i$.

Condition 4 c(i) < c(z) < ... < c(r)

In English, the cows are stepped.

